

CS61A Lecture 42

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Announcements



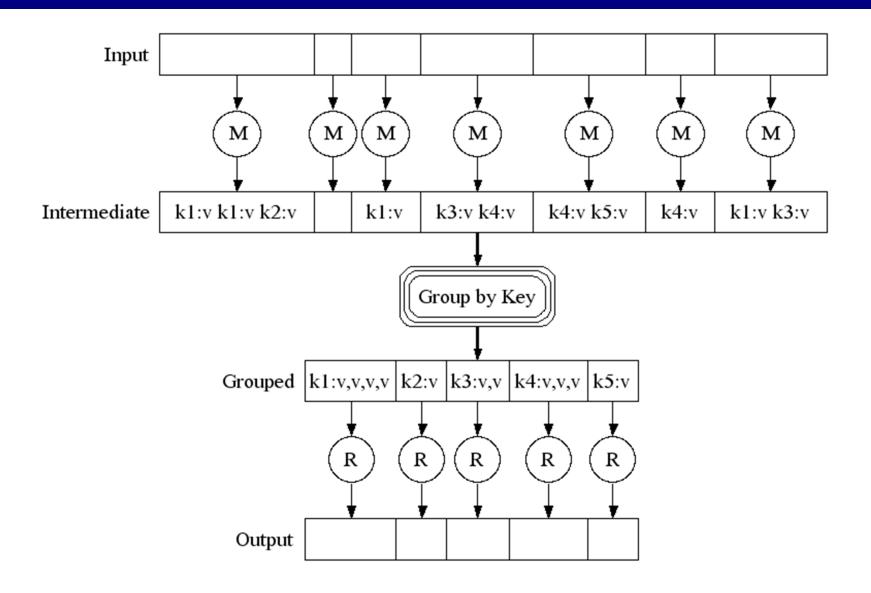
☐ HW13 due Wednesday

□ Scheme project due tonight!!!

□ Scheme contest deadline extended to Friday

MapReduce Execution Model





http://research.google.com/archive/mapreduce-osdi04-slides/index-auto-0007.html





The mapper and reducer are both self-contained Python programs



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Read from standard input and write to standard output!



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Mapper



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Mapper

```
def emit_vowels(line):
    for vowel in 'aeiou':
        count = line.count(vowel)
        if count > 0:
        emit(vowel, count)
```



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Read from standard input and write to standard output!

Mapper

```
#!/usr/bin/env python3
import sys
from ucb import main
from mapreduce import emit

def emit_vowels(line):
    for vowel in 'aeiou':
        count = line.count(vowel)
        if count > 0:
        emit(vowel, count)
```



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Tell Unix: this is Python

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Read from standard input and write to standard output!

Mapper Tell Unix: this is Python #!/usr/bin/env python3 The **emit** function outputs a key import sys and value as a line of text to from ucb import main standard output from mapreduce import emit def emit vowels(line): for vowel in 'aeiou': count = line.count(vowel) if count > 0: emit(vowel, count) for line in sys.stdin: emit vowels(line)



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Mapper Tell Unix: this is Python #!/usr/bin/env python3 The **emit** function outputs a key import sys and value as a line of text to from ucb import main standard output from mapreduce import emit def emit vowels(line): for vowel in 'aeiou': count = line.count(vowel) if count > 0: emit(vowel, count) for line in sys.stdin: Mapper inputs are lines of text emit_vowels(line) provided to standard input



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Reducer



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Read from standard input and write to standard output!

Reducer

```
#!/usr/bin/env python3
import sys
from ucb import main
from mapreduce import emit, group_values_by_key
```



The *mapper* and *reducer* are both self-contained Python programs

Read from standard input and write to standard output!

Reducer

```
#!/usr/bin/env python3
import sys
from ucb import main
from mapreduce import emit, group_values_by_key
Takes and returns iterators
```



The mapper and reducer are both self-contained Python programs

Read from standard input and write to standard output!

Reducer

```
#!/usr/bin/env python3
import sys
from ucb import main
from mapreduce import emit, group_values_by_key
```

Input: lines of text representing key-value pairs, grouped by key

Output: Iterator over (key, value_iterator) pairs that give all values for each key



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Read from standard input and write to standard output!

emit(key, sum(value_iterator))

Reducer

```
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import sys
from ucb import main
from mapreduce import emit, group_values_by_key

Input: lines of text representing key-value pairs, grouped by key

Output: Iterator over (key, value_iterator) pairs that give all values for each key
```

for key, value_iterator in group_values_by_key(sys.stdin):





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- 1. Dense Linear Algebra
- 2. Sparse Linear Algebra
- 3. Spectral Methods
- 4. N-Body Methods
- 5. Sructured Grids
- 6. Unstructured Grids
- 7. MapReduce

- 8. Combinational Logic
- 9. Graph Traversal
- 10. Dynamic Programming
- 11. Backtrack and Branch-and-Bound
- 12. Graphical Models
- 13. Finite State Machines



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MapReduce is only one of these patterns



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MapReduce is only one of these patterns

The rest require shared mutable state





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Processes execute in separate interpreters, generally not sharing data



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Shared state can be communicated explicitly between processes



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The concepts of threads and processes exist in other systems as well

Threads



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```
from threading import Thread, current_thread

def thread_hello():
    other = Thread(target=thread_say_hello, args=())
```



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```
from threading import Thread, current thread
                            Function that the new thread should run
def thread hello():
    other = Thread(target=thread say hello, args=())
    other.start()_
                     Start the other thread
                                              Arguments to
    thread say hello()
def thread say hello():
    print('hello from', current thread().name)
>>> thread hello()
hello from Thread-1
hello from MainThread
```



The **threading** module contains classes that enable threads to be created and synchronized

```
from threading import Thread, current thread
                             Function that the new thread should run
def thread hello():
    other = Thread(target=thread_say_hello, args=())
                     Start the other thread
    other.start()_
                                               Arguments to
    thread say hello()
                                               that function
def thread say hello():
    print('hello from', current thread().name)
>>> thread hello()
                               Print output is not synchronized,
hello from Thread-1
                                 so can appear in any order
hello from MainThread
```

Processes



The **multiprocessing** module contains classes that enable processes to be created and synchronized

Here is a "hello world" example with two processes:

```
from multiprocessing import Process, current_process
                             Function that the new process should run
def process hello():
    other = Process(target=process_say_hello, args=())
                     Start the other process
    other.start()_
                                                 Arguments to
    process say hello()
                                                  that function
def process say hello():
    print('hello from', current process().name)
>>> process hello()
                              Print output is not synchronized,
hello from MainProcess
                                 so can appear in any order
>>> hello from Process-1
```





Shared state that is mutated and accessed concurrently by multiple threads can cause subtle bugs



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Here is an example with two threads that concurrently update a counter:

from threading import Thread



Shared state that is mutated and accessed concurrently by multiple threads can cause subtle bugs

```
from threading import Thread
counter = [0]
```



Shared state that is mutated and accessed concurrently by multiple threads can cause subtle bugs

```
from threading import Thread
counter = [0]
def increment():
```



Shared state that is mutated and accessed concurrently by multiple threads can cause subtle bugs

```
from threading import Thread
counter = [0]
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    counter[0] = counter[0] + 1
```



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other = Thread(target=increment, args=())
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```



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other.start()
increment()
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Shared state that is mutated and accessed concurrently by multiple threads can cause subtle bugs

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from threading import Thread

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def increment():
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other = Thread(target=increment, args=())
other.start()
increment()
other.join()
```



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Here is an example with two threads that concurrently update a counter:

What is the value of counter[0] at the end?



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from threading import Thread

counter = [0]

def increment():
    counter[0] = counter[0] + 1

other = Thread(target=increment, args=())
other.start()
increment()
other.join()
print('count is now', counter[0])

What is the value of counter[0] at the end?
```



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What is the value of counter[0] at the end?

Only the most basic operations in CPython are atomic, meaning that they have the effect of occurring instantaneously



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What is the value of counter[0] at the end?

Only the most basic operations in CPython are atomic, meaning that they have the effect of occurring instantaneously

The counter increment is three basic operations: read the old value, add 1 to it, write the new value





We can see what happens if a switch occurs at the wrong time by trying to force one in CPython:



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increment()
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We can see what happens if a switch occurs at the wrong time by trying to force one in CPython:

```
from threading import Thread
from time import sleep
counter = [0]
def increment():
    count = counter[0]
    counter[0] = count + 1
other = Thread(target=increment, args=())
other.start()
increment()
other.join()
print('count is now', counter[0])
```



We can see what happens if a switch occurs at the wrong time by trying to force one in CPython:

```
from threading import Thread
from time import sleep
counter = [0]
def increment():
    count = counter[0]
    sleep(0)
    counter[0] = count + 1
other = Thread(target=increment, args=())
other.start()
increment()
other.join()
print('count is now', counter[0])
```



We can see what happens if a switch occurs at the wrong time by trying to force one in CPython:

```
from threading import Thread
from time import sleep
counter = [0]
def increment():
    count = counter[0]
    sleep(0) — May cause the interpreter to switch threads
    counter[0] = count + 1
other = Thread(target=increment, args=())
other.start()
increment()
other.join()
print('count is now', counter[0])
```







Given a switch at the **sleep** call, here is a possible sequence of operations on each thread:

Thread 0

Thread 1



```
Thread 0 Thread 1 read counter[0]: 0
```



```
Thread 0 Thread 1 read counter[0]: 0 read counter[0]: 0
```



```
Thread 0 Thread 1
read counter[0]: 0
read counter[0]: 0
calculate 0 + 1: 1
```









Given a switch at the **sleep** call, here is a possible sequence of operations on each thread:

The counter ends up with a value of 1, even though it was incremented twice!





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Managing shared state is a key challenge in parallel computing

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Managing shared state is a key challenge in parallel computing

- Under-synchronization doesn't protect against race conditions and other parallel bugs
- Over-synchronization prevents non-conflicting accesses from occurring in parallel, reducing a program's efficiency



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We will see some basic tools for managing shared state







```
from queue import Queue
queue = Queue()
def increment():
    count = queue.get()
    sleep(0)
    queue.put(count + 1)
other = Thread(target=increment, args=())
other.start()
queue.put(0)
increment()
other.join()
print('count is now', queue.get())
```



```
from queue import Queue
                                Synchronized FIFO queue
queue = Queue()
def increment():
    count = queue.get()
    sleep(0)
    queue.put(count + 1)
other = Thread(target=increment, args=())
other.start()
queue.put(0)
increment()
other.join()
print('count is now', queue.get())
```



```
from queue import Queue
                                 Synchronized FIFO queue
queue = Queue()
def increment():
                               Waits until an item is available
    count = queue.get() =
    sleep(0)
    queue.put(count + 1)
other = Thread(target=increment, args=())
other.start()
queue.put(0)
increment()
other.join()
print('count is now', queue.get())
```



```
from queue import Queue
                                  Synchronized FIFO queue
queue = Queue()
def increment():
                                Waits until an item is available
    count = queue.get() =
    sleep(0)
    queue.put(count + 1)
other = Thread(target=increment, args=())
other.start()
queue.put(0)
                       Add initial value of 0
increment()
other.join()
print('count is now', queue.get())
```