

# CS61A Lecture 35

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#### Announcements



- ☐ HW11 due next Wednesday
- ☐ Scheme project out

# The way in which names are looked up in Scheme and Python is called *lexical scope* (or *static scope*) Lexical scope: The parent of a frame is the environment in which a procedure was *defined*Dynamic scope: The parent of a frame is the environment in which a procedure was *called*Special form to create dynamically scoped procedures

**Dynamic Scope** 

 $\mbox{(g 3 7)} \label{eq:gamma}$  Lexical scope: The parent for £'s frame is the global frame

(define g (lambda (x y) (f (+ x x))))

(x) (+ x y)))

Error: unknown identifier: y **Dynamic scope**: The parent for £'s frame is g's frame

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# **Functional Programming**



All functions are pure functions

No re-assignment and no mutable data types

Name-value bindings are permanent

Advantages of functional programming:

- The value of an expression is independent of the order in which subexpressions are evaluated
- Sub-expressions can safely be evaluated in parallel or lazily
- Referential transparency: The value of an expression does not change when we substitute one of its sub-expression with the value of that sub-expression

But... Can we make basic loops efficient?

Yes!

## Iteration Versus Recursion in Python



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In Python, recursive calls always create new active frames

#### **Iteration and Recursion**



Reminder: Iteration is a special case of recursion

Idea: The state of iteration can be passed as parameters

```
def factorial(n):
    total = 1
    while n > 0:
        n, total = n - 1, total * n
    return total

def factorial(n, total):
    if n == 0:
        return total
    return factorial(n - 1, total * n)
```

But this converted version still uses linear space in Python

## **Tail Recursion**



From the Revised<sup>7</sup> Report on the Algorithmic Language Scheme:

"Implementations of Scheme are required to be **properly tail-recursive**. This allows the execution of an iterative computation in constant space, even if the iterative computation is described by a syntactically recursive procedure."

## Tail Calls



A procedure call that has not yet returned is *active*. Some procedure calls are *tail calls*. A Scheme interpreter should support an unbounded number of active tail calls.

A tail call is a call expression in a tail context:

- The last body sub-expression in a lambda expression
- Sub-expressions 2 & 3 in a tail context if expression
- All non-predicate sub-expressions in a tail context cond
- The last sub-expression in a tail context and or or
- The last sub-expression in a tail context begin

# Example: Length of a List



A call expression is not a tail call if more computation is still required in the calling procedure.

Linear recursions can often be rewritten to use tail calls

## **Eval with Tail Call Optimization**



The return value of the tail call is the return value of the current procedure call.

Therefore, tail calls shouldn't increase the environment size.

In the interpreter, recursive calls to **scheme\_eval** for tail calls must instead be expressed iteratively.

## Logical Special Forms, Revisited



Logical forms may only evaluate some sub-expressions.

- And and or: (and <e<sub>1</sub>> ... <e<sub>n</sub>>), (or <e<sub>1</sub>> ... <e<sub>n</sub>>)
- Cond expr'n: (cond (<p<sub>1</sub>> <e<sub>1</sub>>) ... (<p<sub>n</sub>> <e<sub>n</sub>>) (else <e>))

The value of an **if** expression is the value of a sub-expression.

- Evaluate the predicate.

   Choose a sub-expression: <consequent> or <alternative>
- Evaluate that sub-expression in place of the whole expression.

scheme\_eval )

Evaluation of the tail context does not require a recursive call.

E.g., replace (if false 1 (+ 2 3)) with (+ 2 3) and iterate.

## Example: Reduce



```
(define (reduce procedure s start)

(if (null? s) start

(reduce procedure

(cdr s)

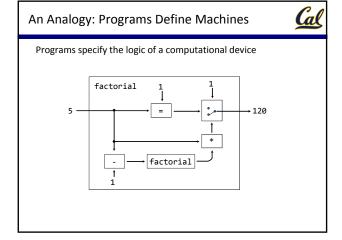
(procedure start (car s)))))

Recursive call is a tail call.

Other calls are not; constant space depends on procedure.
```

```
(reduce * '(3 4 5) 2) 120
(reduce (lambda (x y) (cons y x)) '(3 4 5) '(2)) (5 4 3 2)
```

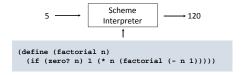
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## Interpreters are General Computing Machines



An interpreter can be parameterized to simulate any machine



Our Scheme interpreter is a universal machine

A bridge between the data objects that are manipulated by our programming language and the programming language itself Internally, it is just a set of manipulation rules

## Interpretation in Python



eval: Evaluates an expression in the current environment and returns the result. Doing so may affect the environment.

**exec**: Executes a statement in the current environment. Doing so may affect the environment.

os.system('python <file>'): Directs the operating
system to invoke a new instance of the Python interpreter.