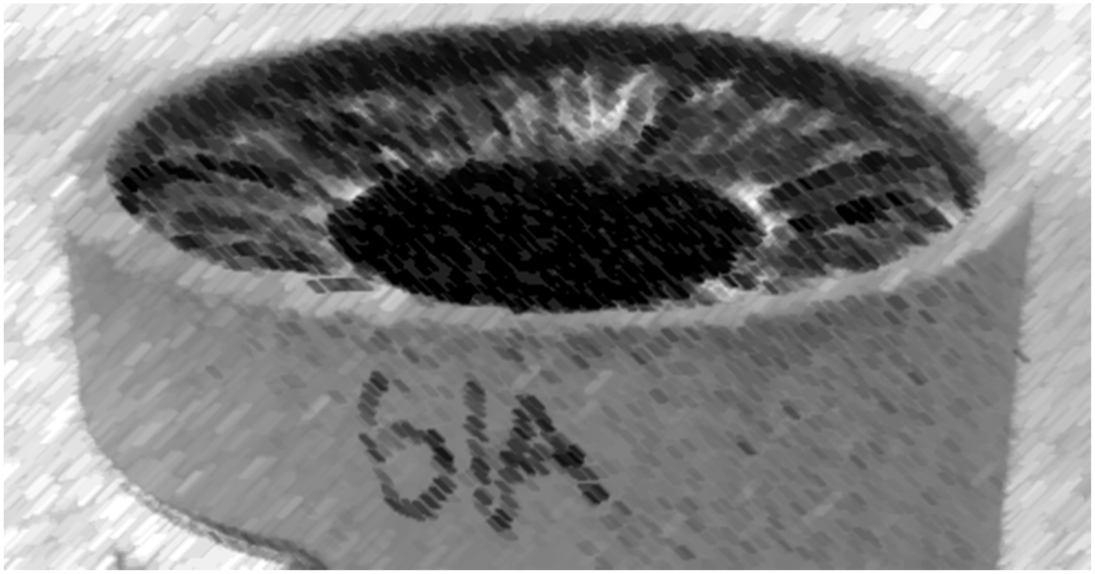

CS61A Lecture 30

Amir Kamil

UC Berkeley

April 1, 2013

Announcements

- HW9 due Wednesday
- Ants extra credit due Wednesday
  - See Piazza for submission instructions
- Hog revisions out, due next Monday

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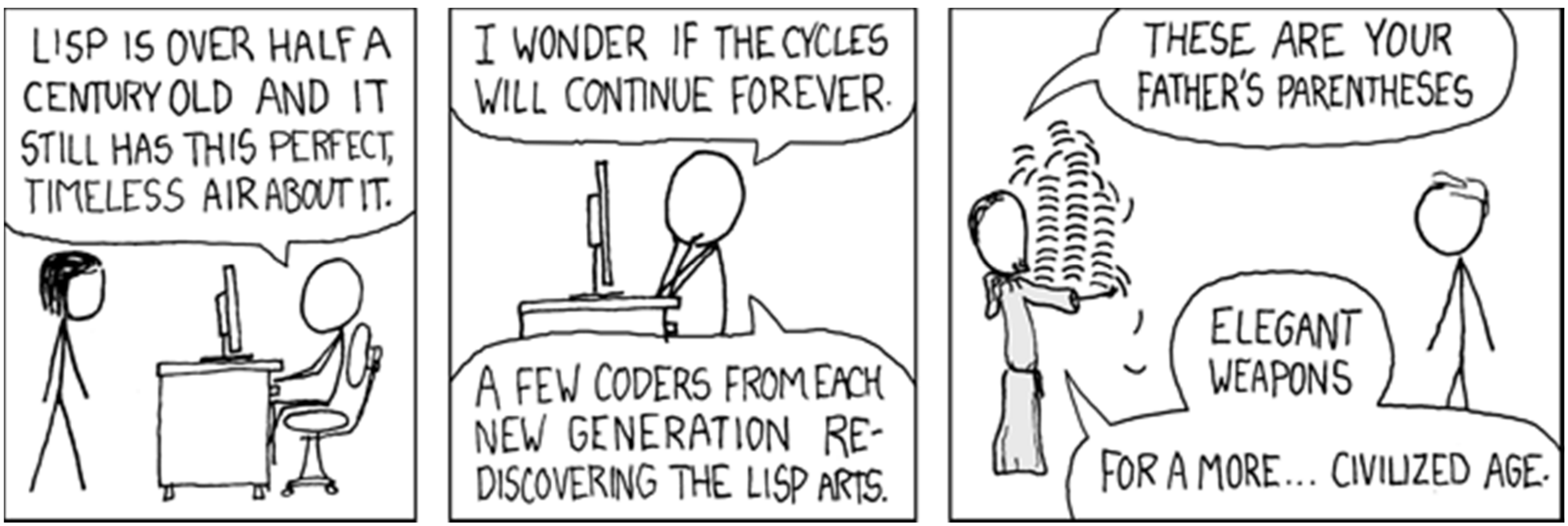

http://imgs.xkcd.com/comics/lisp_cycles.png

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Evaluates to the *add-x-&-y-&-z*^{2} procedure

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> (define x (cons 1 (cons 2 (cons 3 (cons 4 nil)))))
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A recursive list can be represented as a pair in which the second element is a recursive list or the empty list

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Not a well-formed list!

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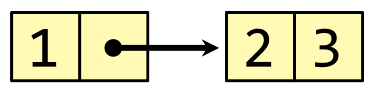

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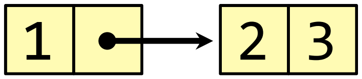

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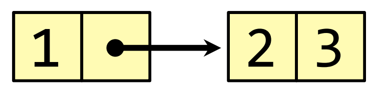

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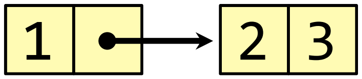

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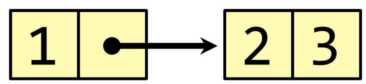

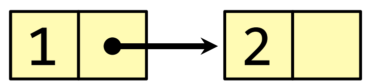

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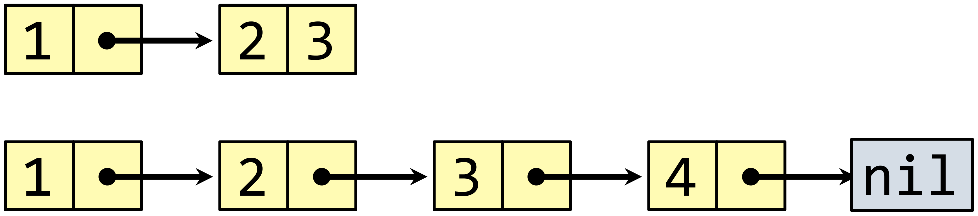

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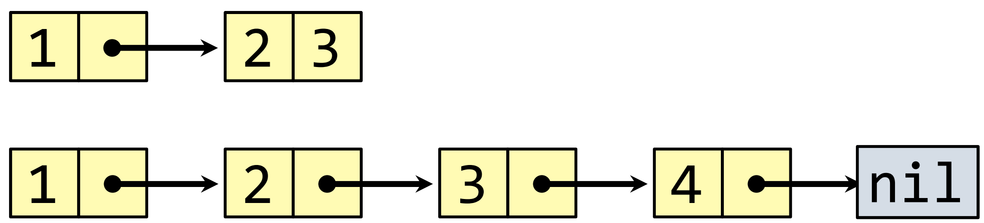

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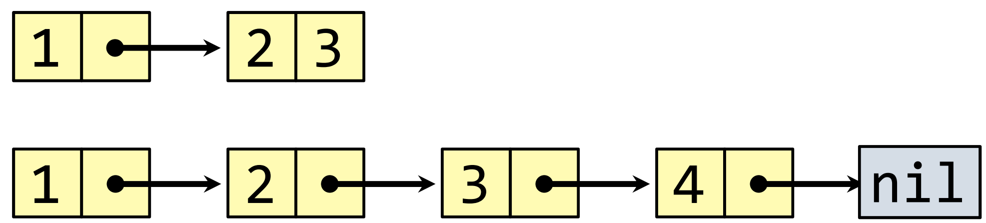

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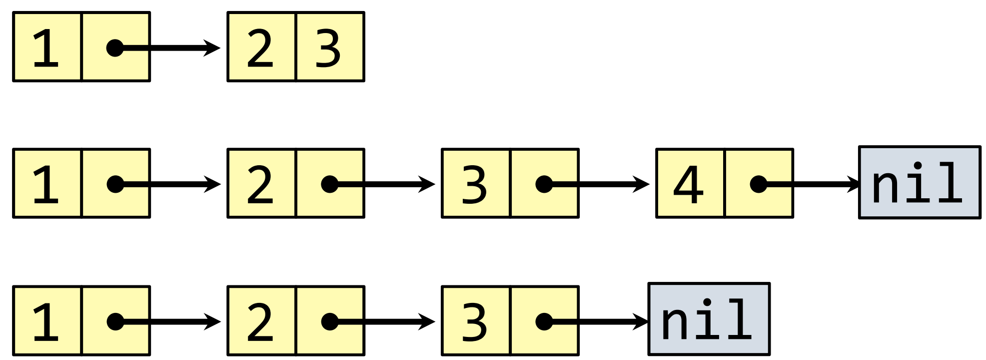

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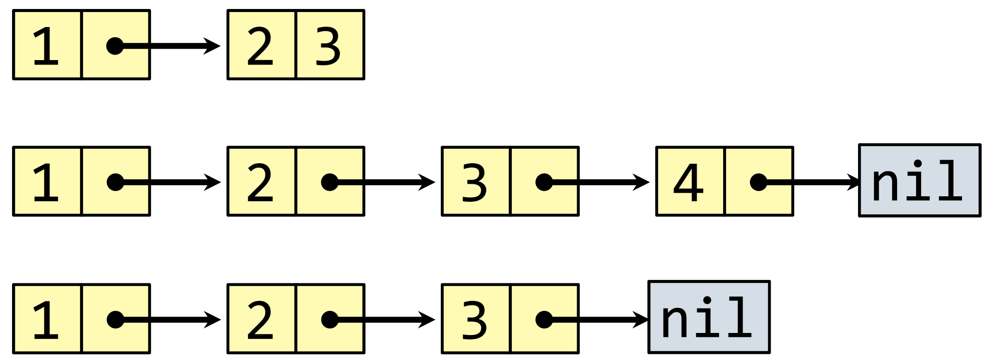

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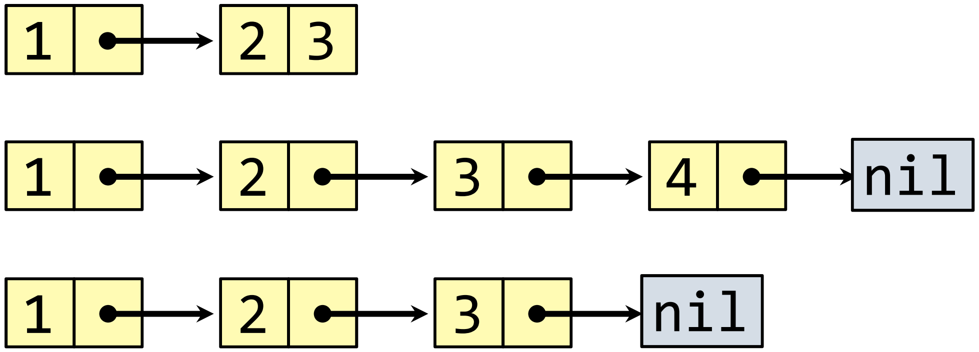

What is the printed result of evaluating this expression?

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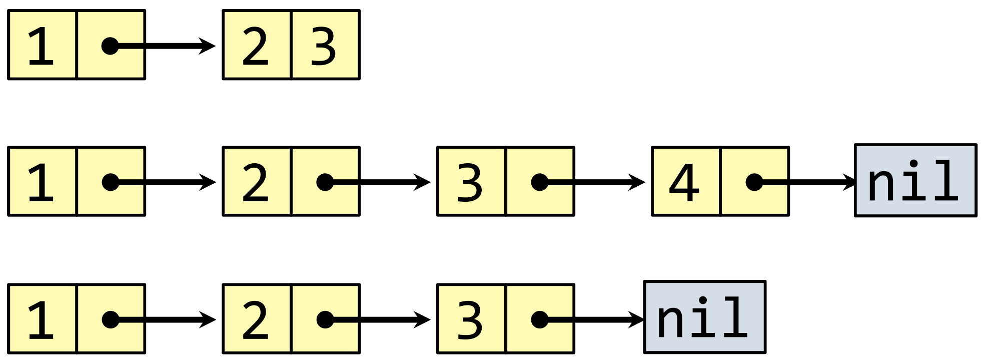

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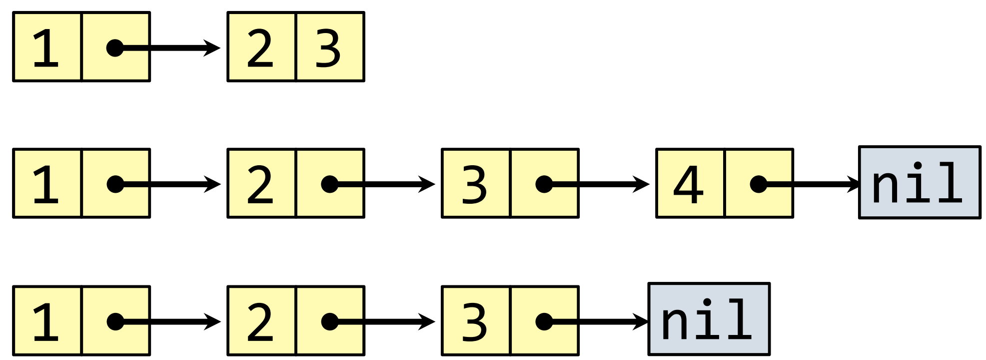

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```