



## CS61A Lecture 26

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UC Berkeley

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# Announcements



- HW9 out tonight, due 4/3
  
- Ants extra credit due 4/3
  - See Piazza for submission instructions

# Data Structure Applications



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The data structures we cover in 61A are used everywhere in CS

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More about data structures in 61B

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Example: recursive lists (also called *linked lists*)

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Example: recursive lists (also called *linked lists*)

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- Interpreters and compilers
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# Data Structure Applications



The data structures we cover in 61A are used everywhere in CS

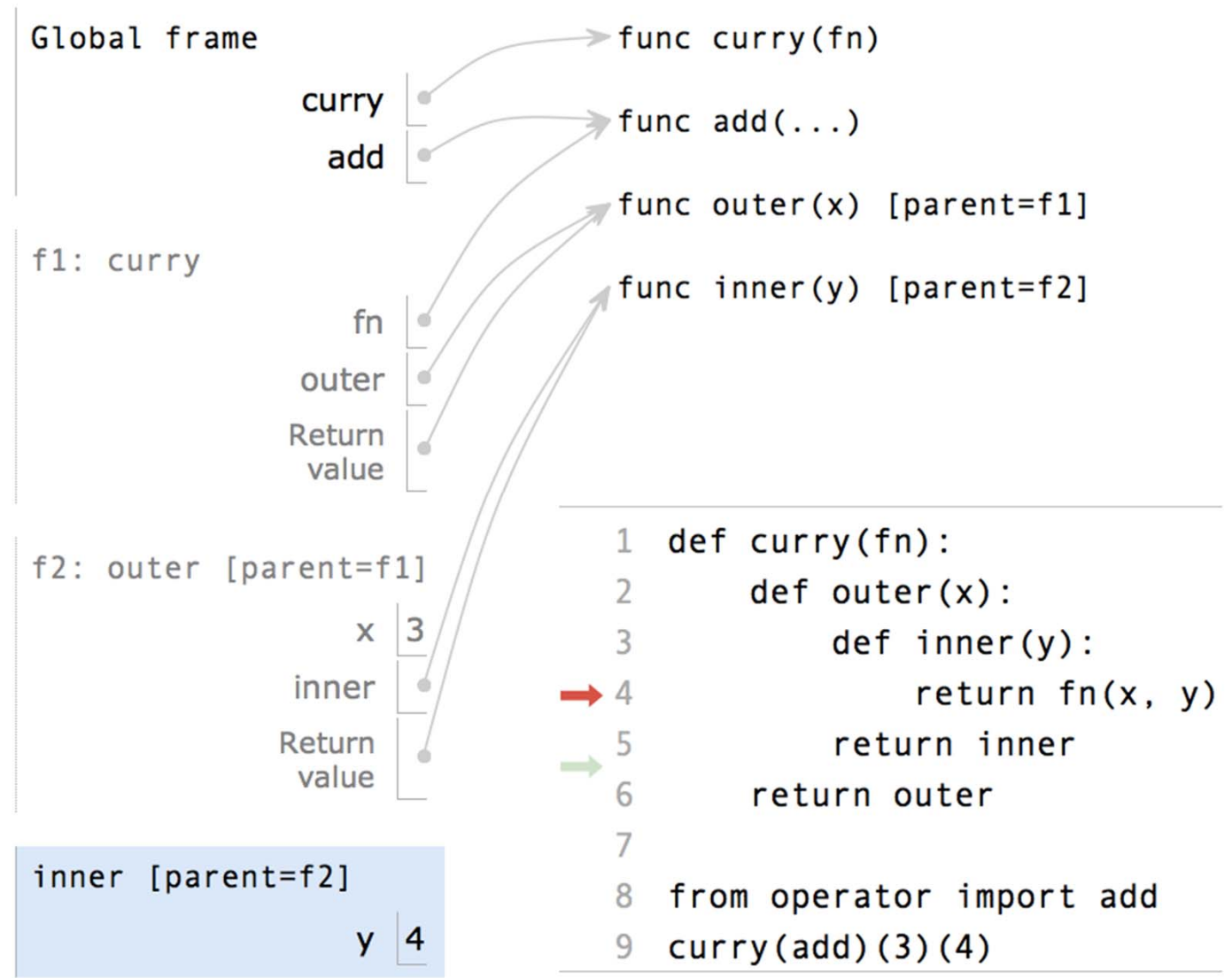
More about data structures in 61B

Example: recursive lists (also called *linked lists*)

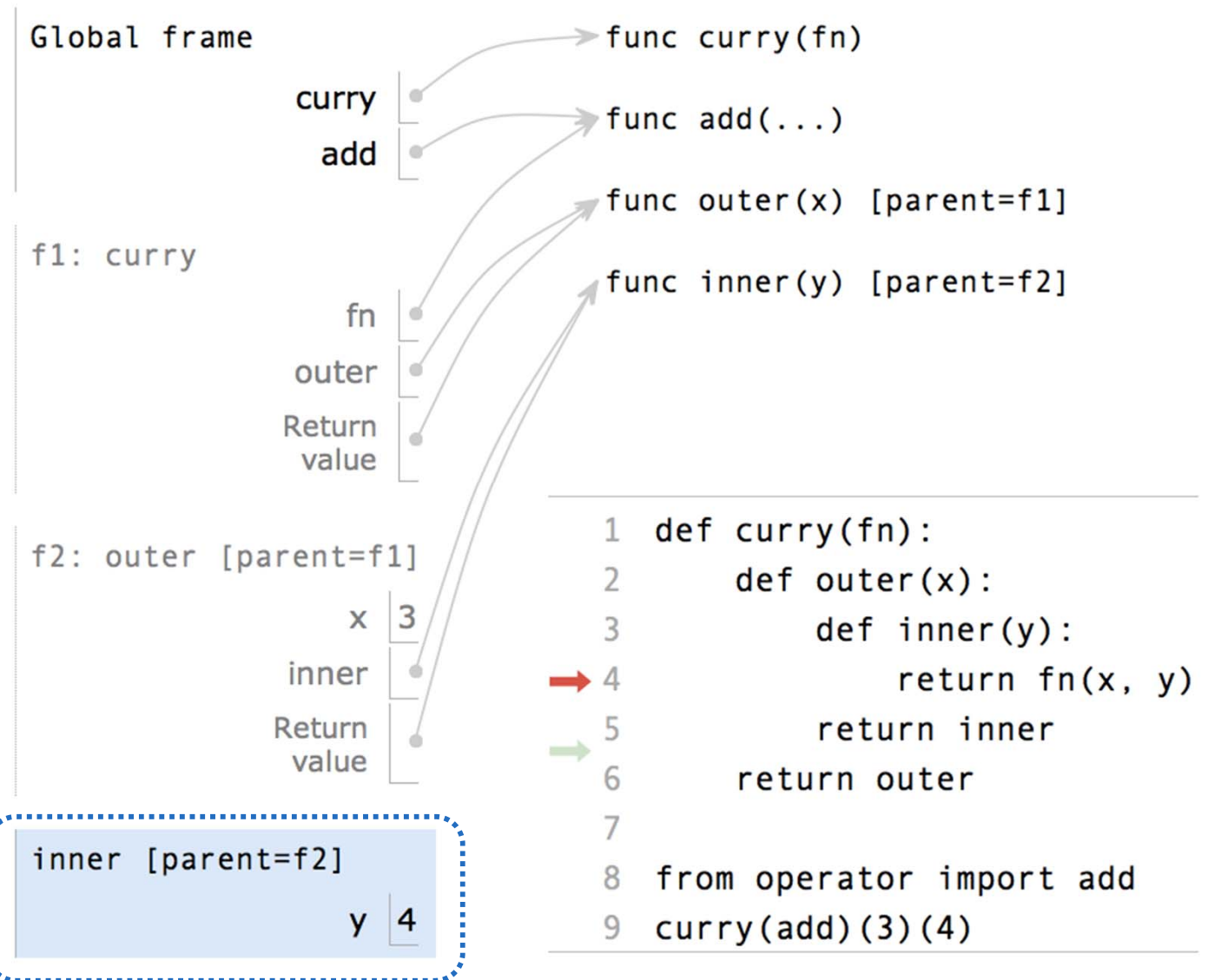
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The Scheme programming language, which we will learn soon, uses recursive lists as its primary data structure

# Example: Environments

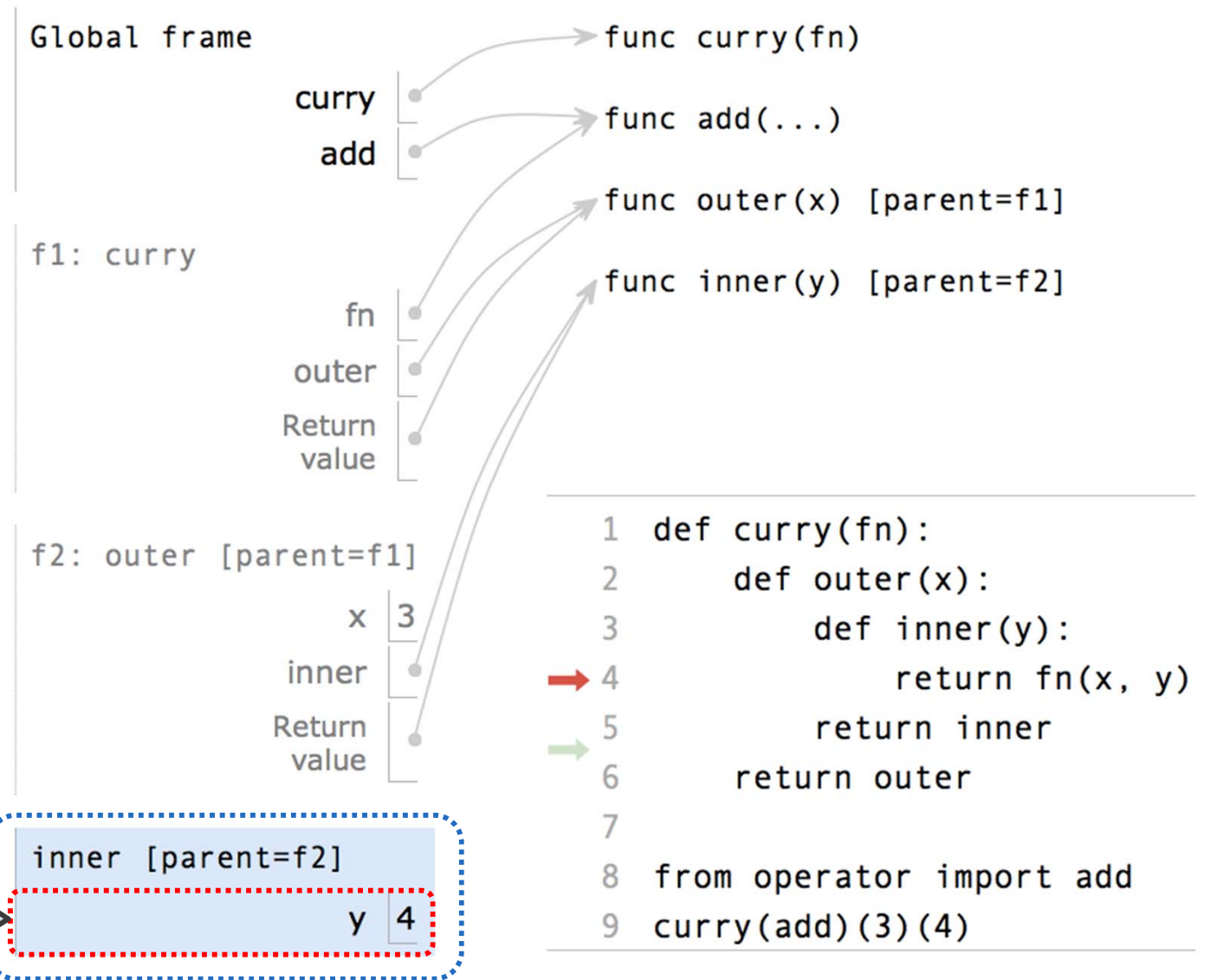


# Example: Environments



Recursive List

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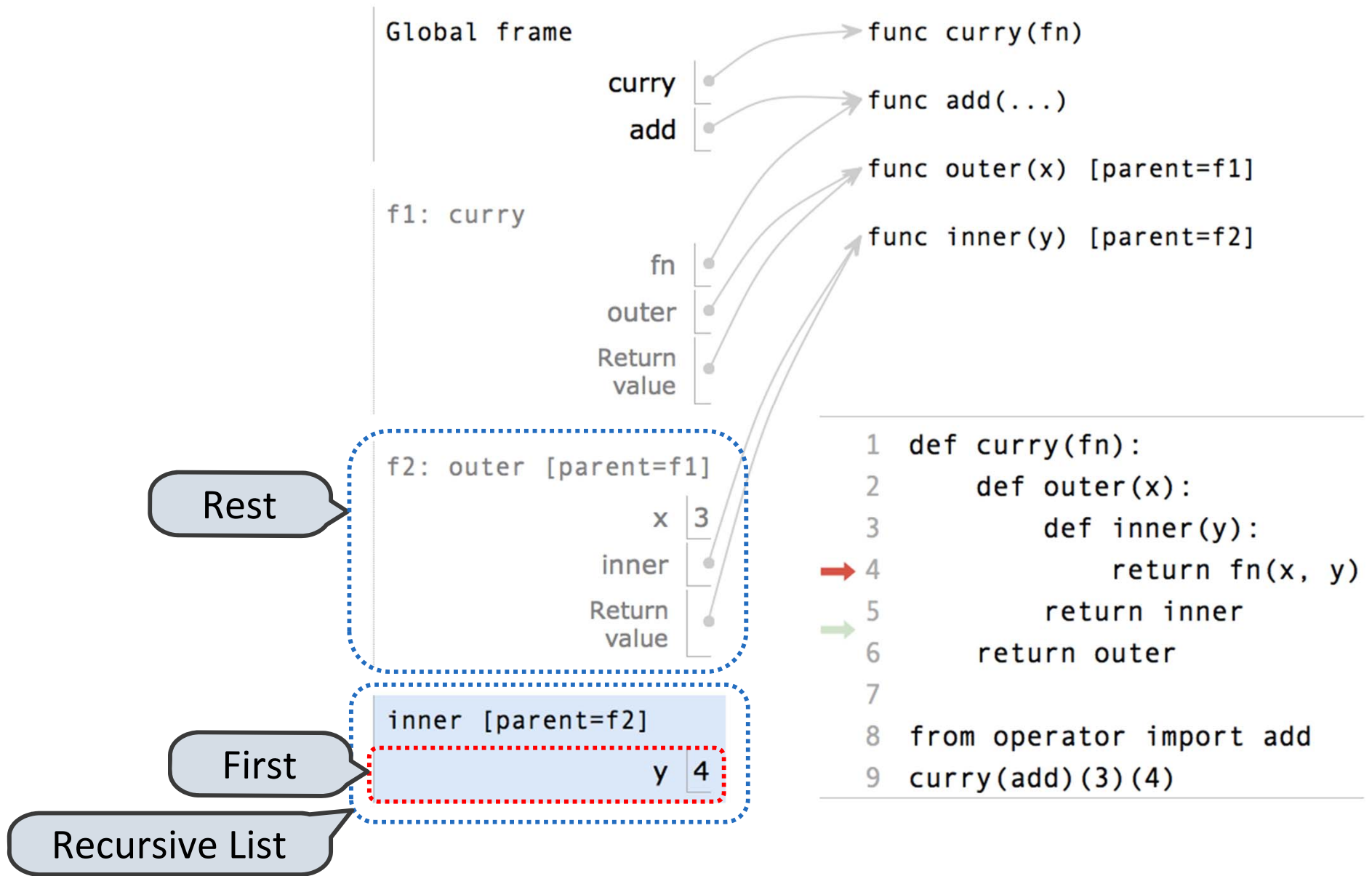


First

Recursive List

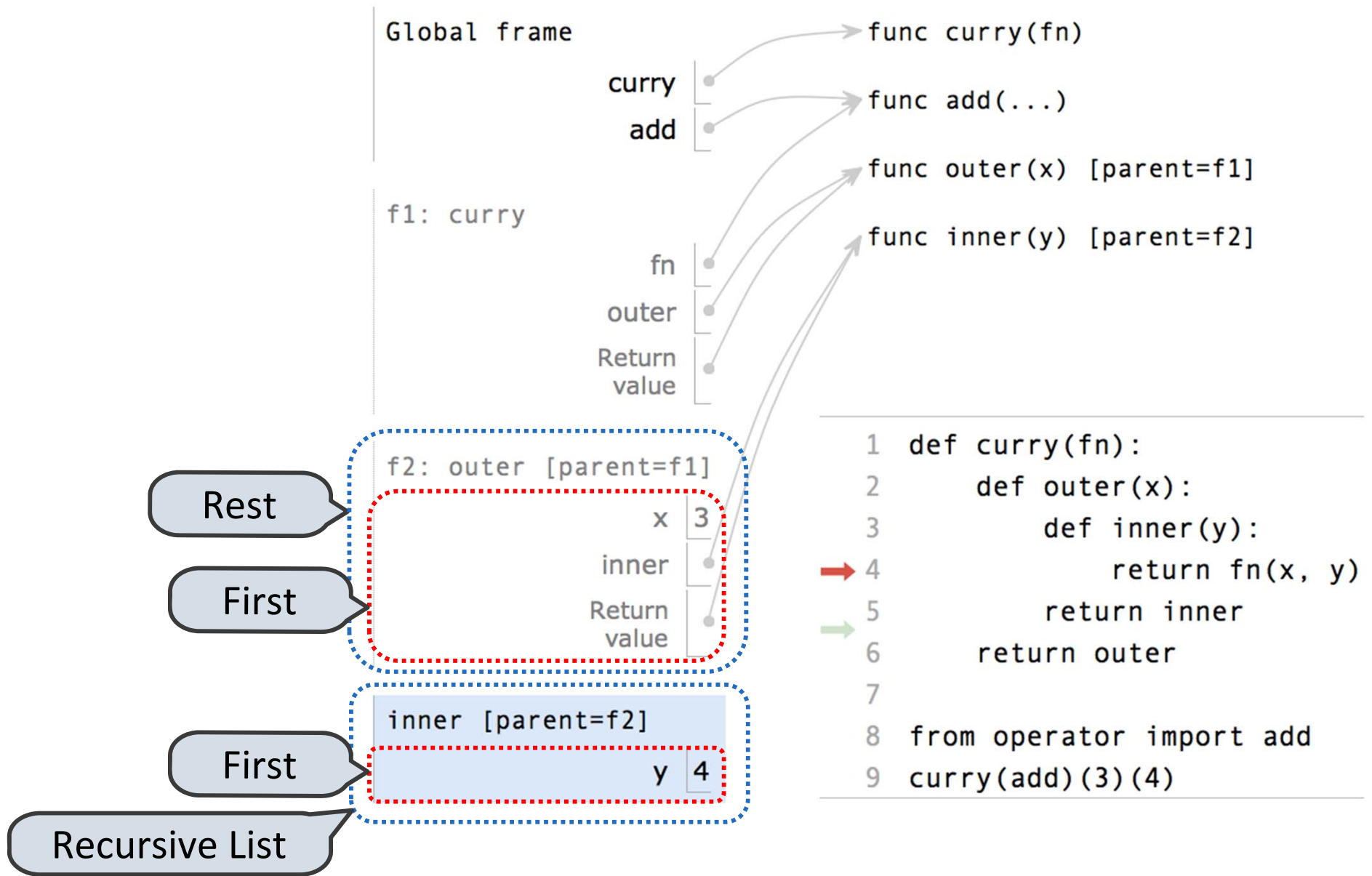
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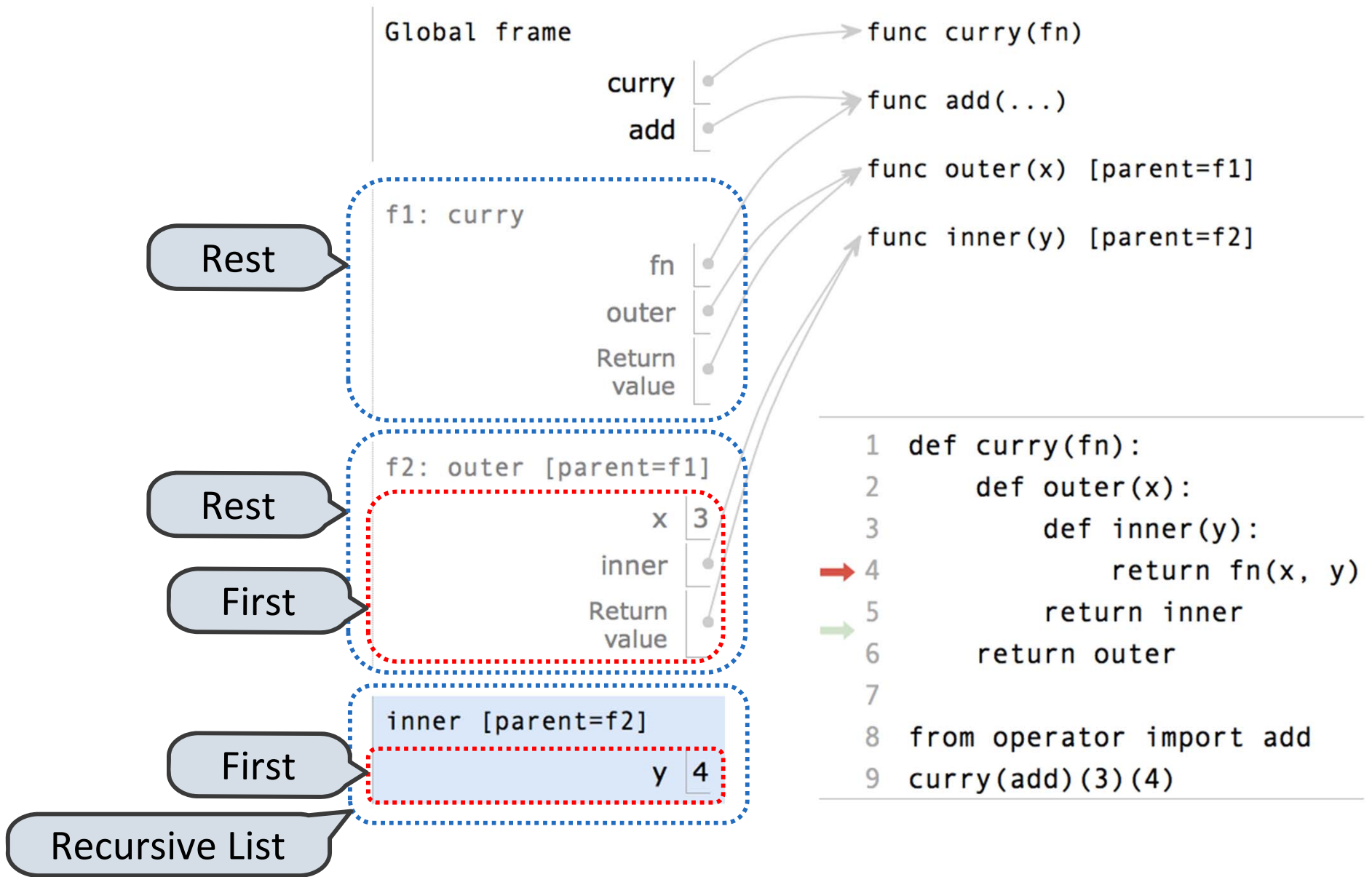
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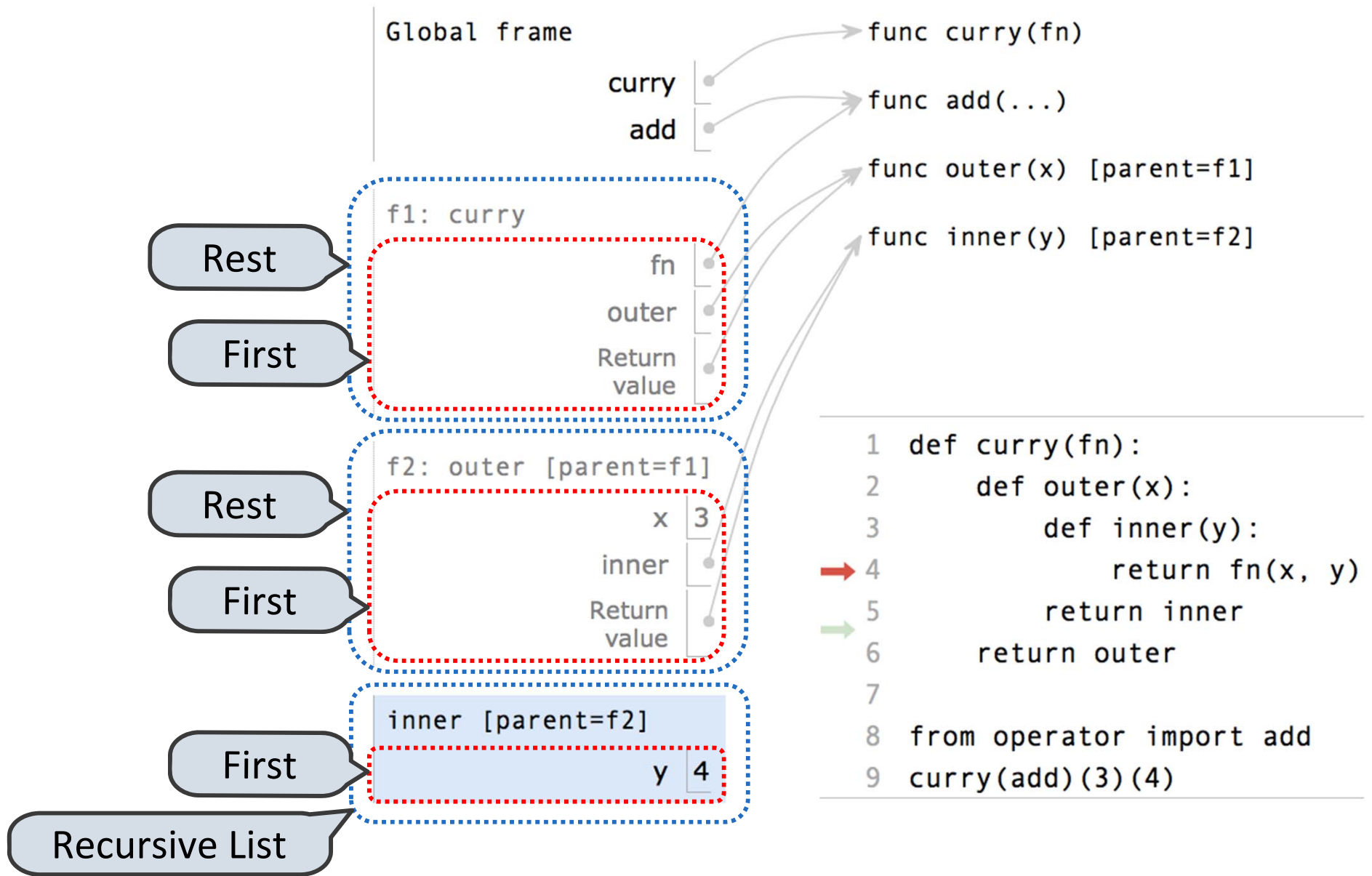
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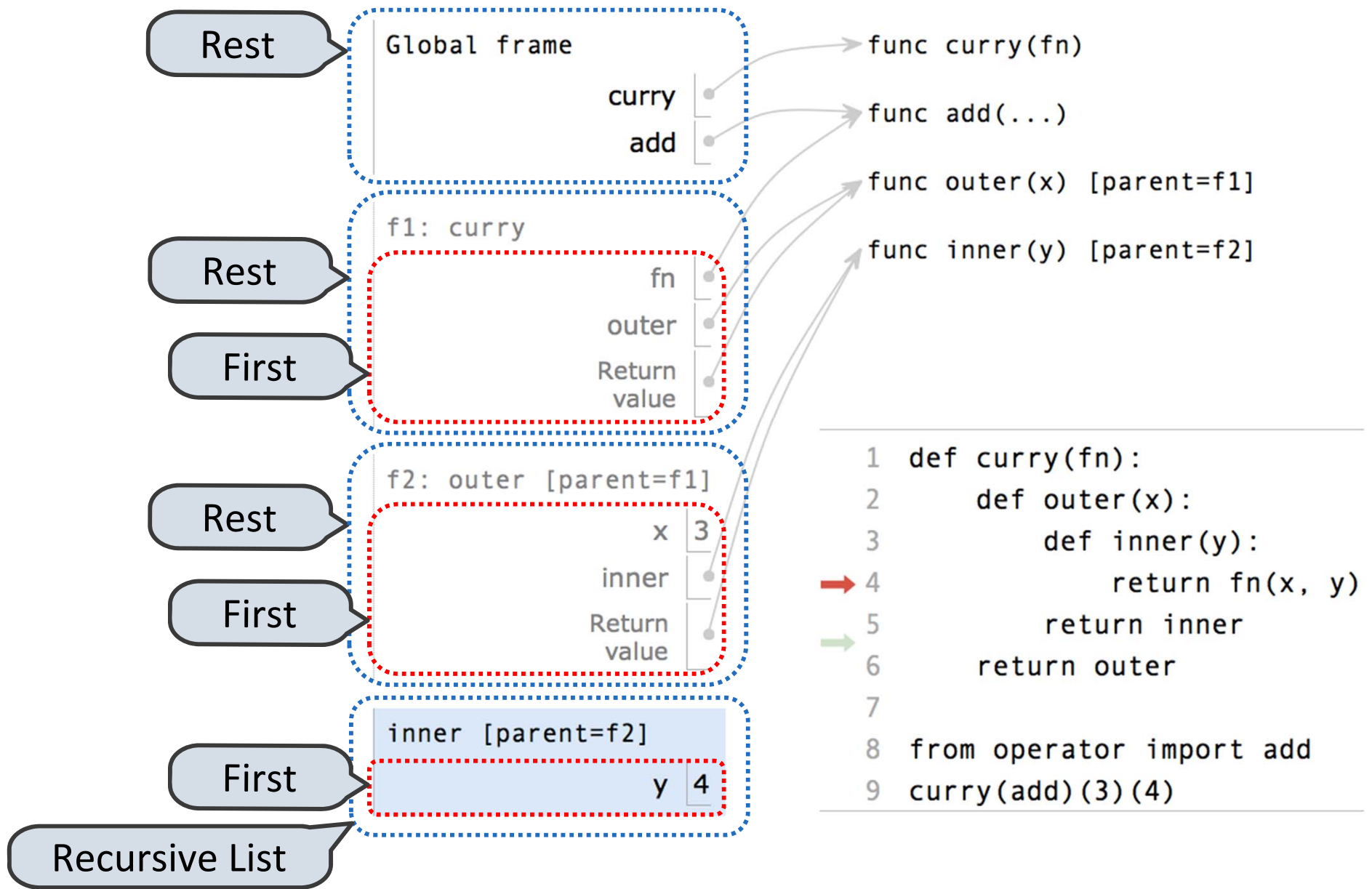


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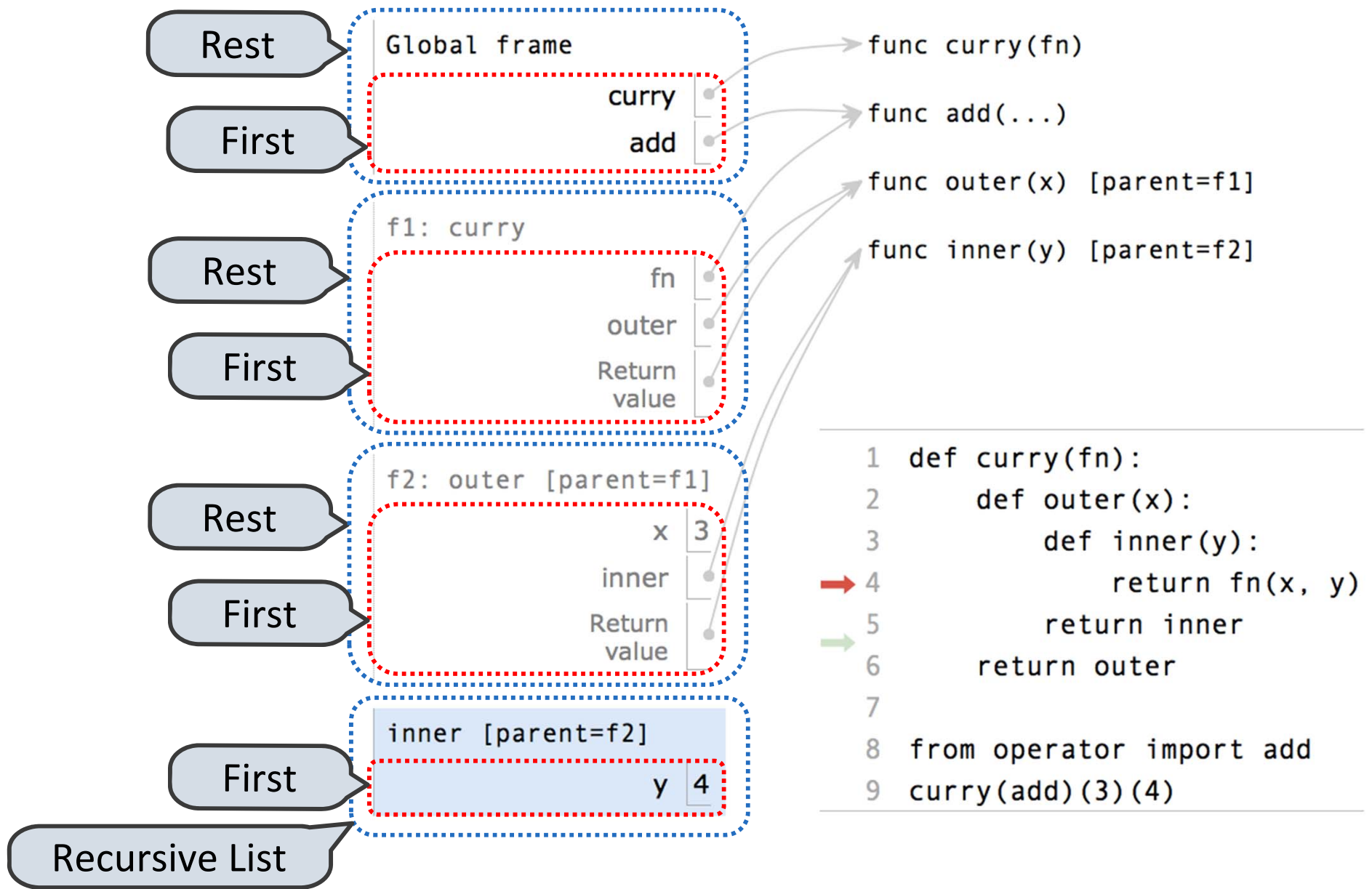
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# Example: Environments



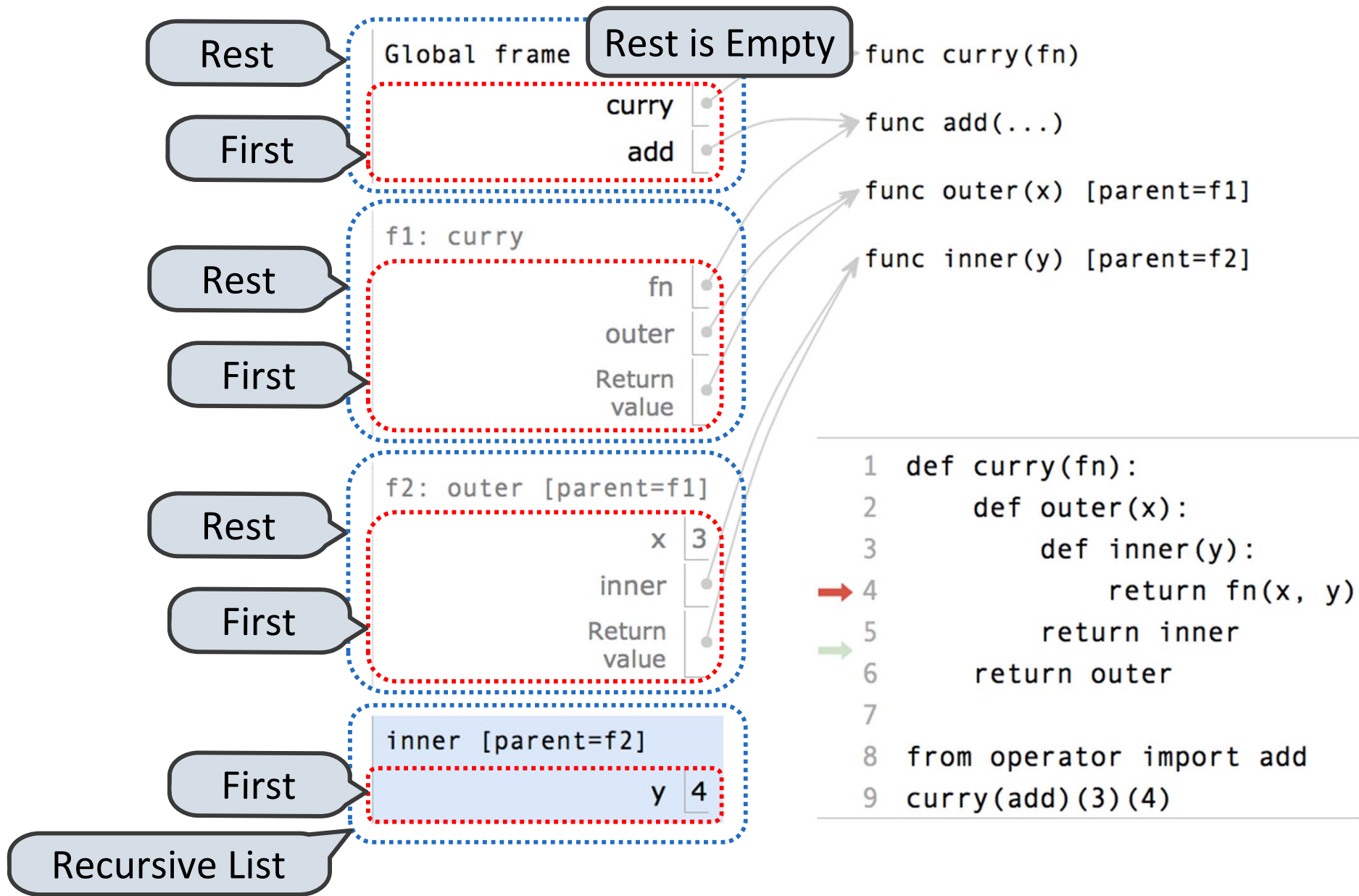
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# Example: Environments



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# Example: Environments



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# Trees with Internal Node Values



Trees can have values at internal nodes as well as their leaves.

```
class Tree(object):
    def __init__(self, entry, left=None, right=None):
        self.entry = entry
        self.left = left
        self.right = right

def fib_tree(n):
    if n == 1:
        return Tree(0)
    if n == 2:
        return Tree(1)
    left = fib_tree(n - 2)
    right = fib_tree(n - 1)
    return Tree(left.entry + right.entry, left, right)
```

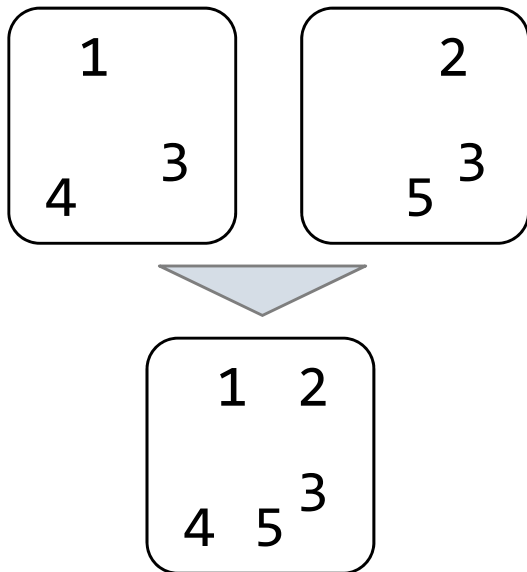
# Implementing Sets



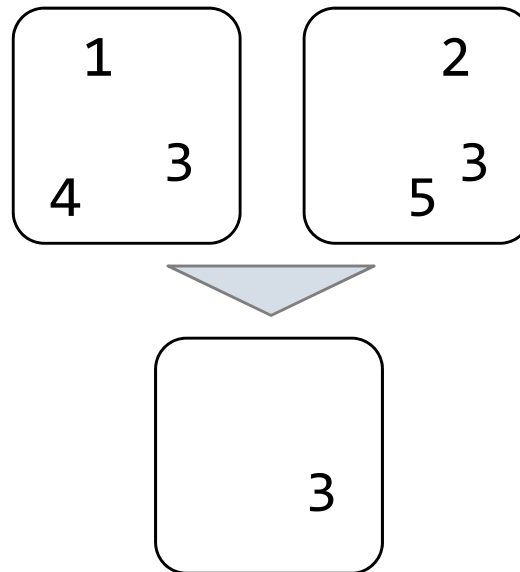
What we should be able to do with a set:

- Membership testing: Is a value an element of a set?
- Union: Return a set with all elements in *set1* **or** *set2*
- Intersection: Return a set with any elements in *set1* **and** *set2*
- Adjunction: Return a set with all elements in *s* and a value *v*

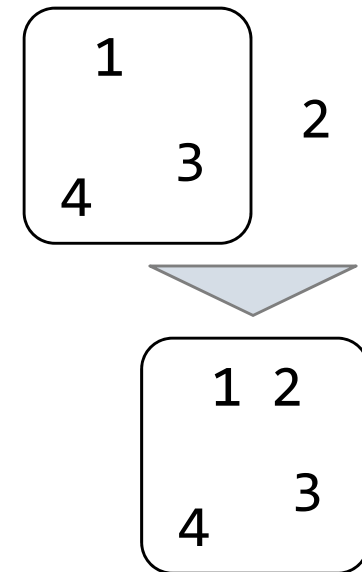
## Union



## Intersection



## Adjunction



# Sets as Unordered Sequences



**Proposal 1:** A set is represented by a recursive list that contains no duplicate items

This is how we implemented dictionaries

```
def empty(s):
    return s is Rlist.empty

def set_contains(s, v):
    if empty(s):
        return False
    elif s.first == v:
        return True
    return set_contains(s.rest, v)
```

# Sets as Unordered Sequences

---





# Sets as Unordered Sequences



```
def adjoin_set(s, v):
```

# Sets as Unordered Sequences



```
def adjoin_set(s, v):  
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# Sets as Unordered Sequences



```
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```

# Sets as Unordered Sequences



```
def adjoin_set(s, v):  
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    return Rlist(v, s)
```

# Sets as Unordered Sequences



Time order of growth

```
def adjoin_set(s, v):  
    if set_contains(s, v):  
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The size of  
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def intersect_set(set1, set2):  
    f = lambda v: set_contains(set2, v)
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```
def intersect_set(set1, set2):  
    f = lambda v: set_contains(set2, v)  
    return filter_rlist(set1, f)
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def union_set(set1, set2):  
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def union_set(set1, set2):  
    f = lambda v: not set_contains(set2, v)  
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    return extend_rlist(set1_not_set2, set2)
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**Proposal 2:** A set is represented by a recursive list with unique elements ordered from least to greatest

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# Set Intersection Using Ordered Sequences

---



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        return Rlist(e1, rest)
    elif e1 < e2:
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    elif e2 < e1:
        return intersect_set2(set1, set2.rest)
```

Order of growth?  $\Theta(n)$

# Tree Sets



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**Proposal 3:** A set is represented as a Tree. Each entry is:

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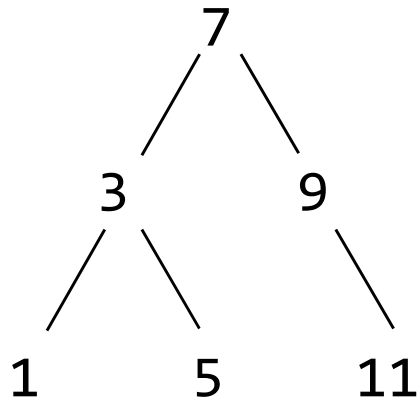
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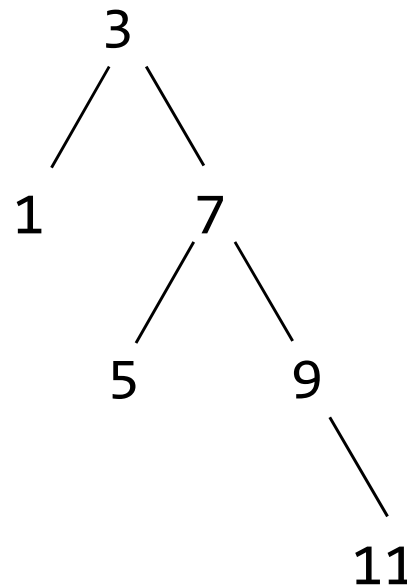
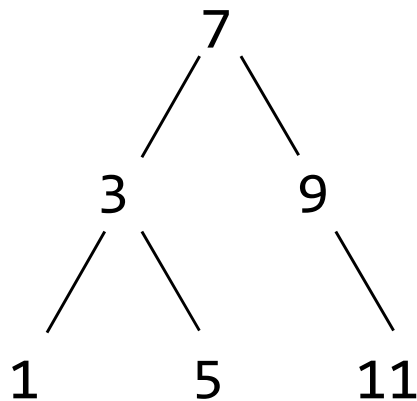


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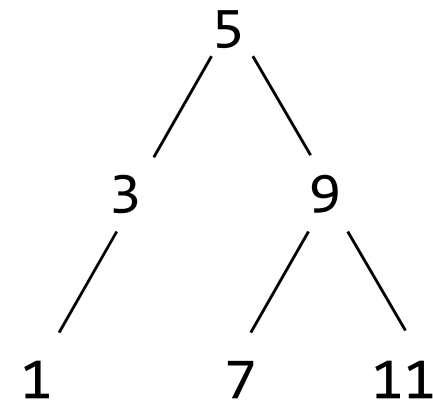
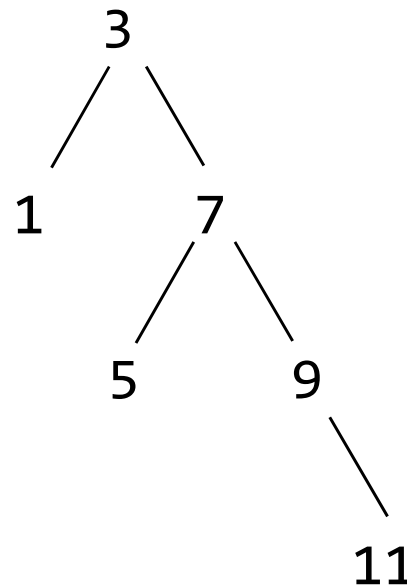
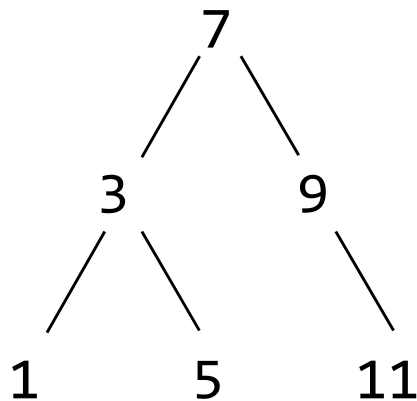


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# Membership in Tree Sets



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Set membership tests traverse the tree

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```

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def set_contains3(s, v):  
    if s is None:
```

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Set membership tests traverse the tree

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def set_contains3(s, v):  
    if s is None:  
        return False
```

# Membership in Tree Sets



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def set_contains3(s, v):  
    if s is None:  
        return False  
    elif s.entry == v:
```



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def set_contains3(s, v):  
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    elif s.entry == v:  
        return True  
    elif s.entry < v:  
        return set_contains3(s.right, v)
```

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        return set_contains3(s.right, v)  
    elif s.entry > v:  
        return set_contains3(s.left, v)
```

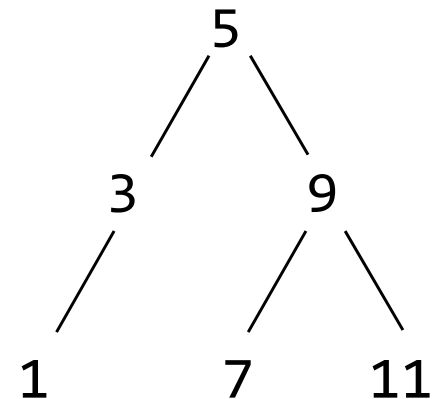
# Membership in Tree Sets



Set membership tests traverse the tree

- The element is either in the left or right sub-branch
- By focusing on one branch, we reduce the set by about half

```
def set_contains3(s, v):  
    if s is None:  
        return False  
    elif s.entry == v:  
        return True  
    elif s.entry < v:  
        return set_contains3(s.right, v)  
    elif s.entry > v:  
        return set_contains3(s.left, v)
```



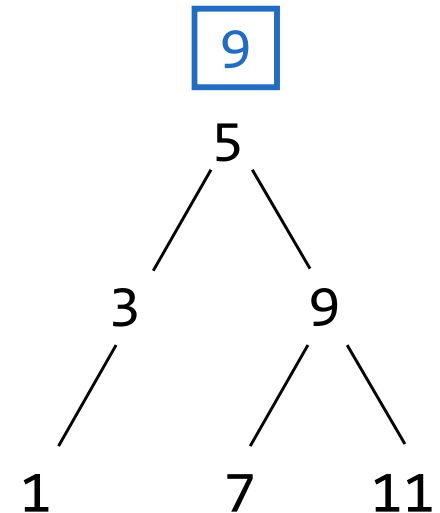
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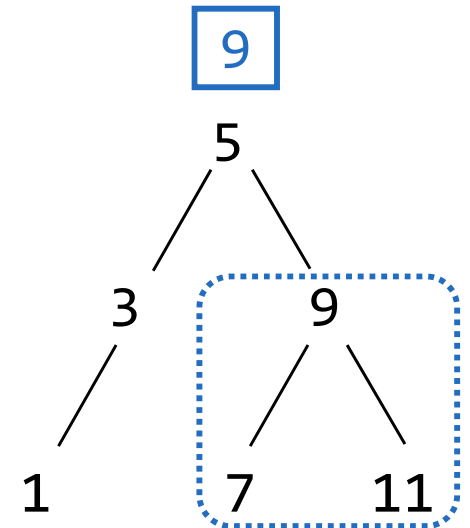
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If 9 is in the set, it is in this branch



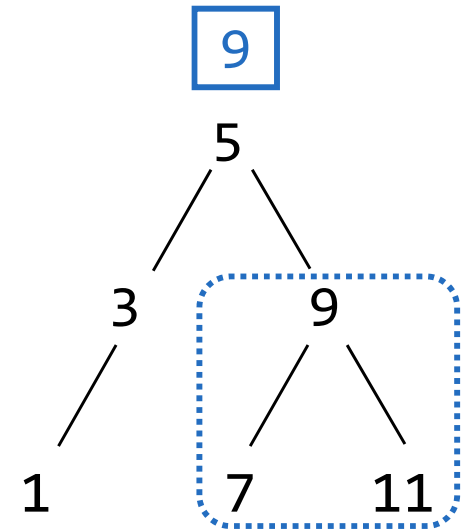
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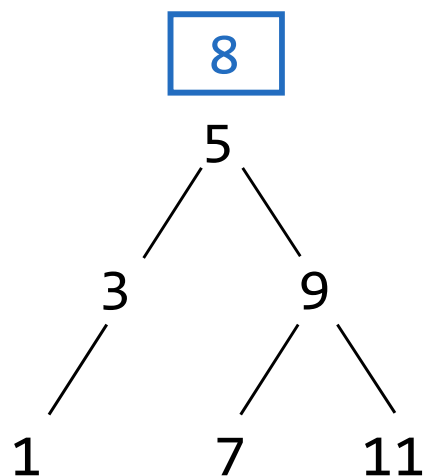
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Order of growth?

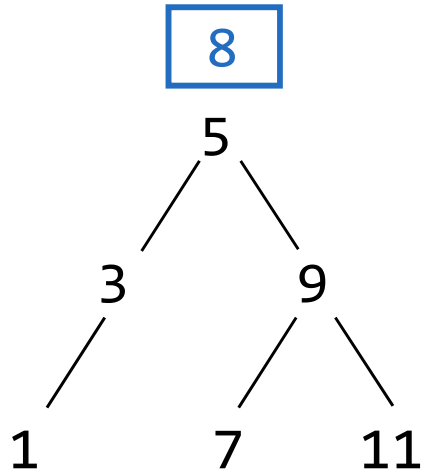
# Adjoining to a Tree Set



# Adjoining to a Tree Set

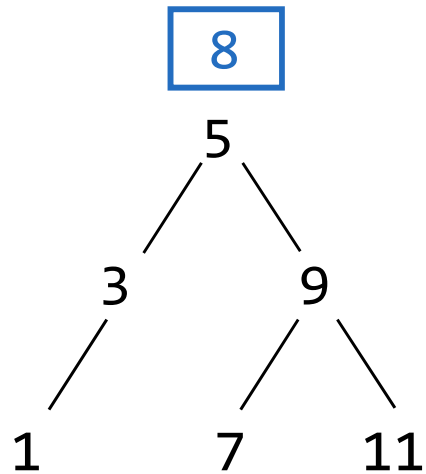


# Adjoining to a Tree Set



Right!

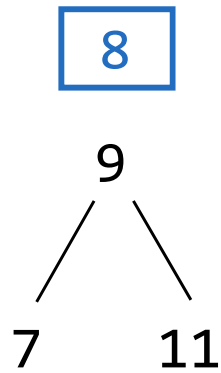
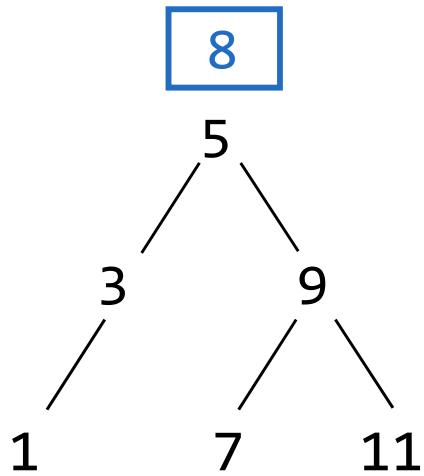
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Right!



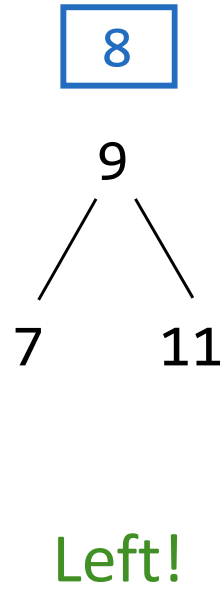
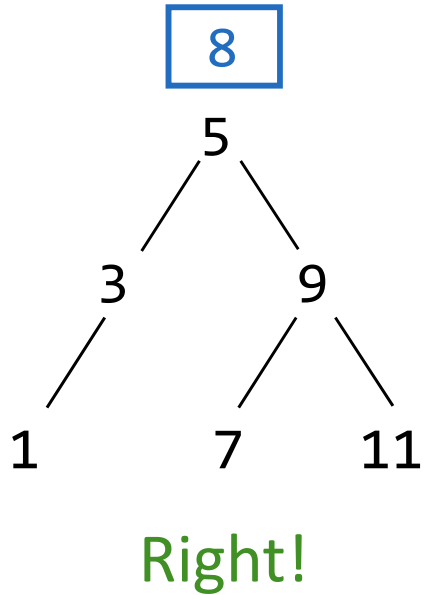
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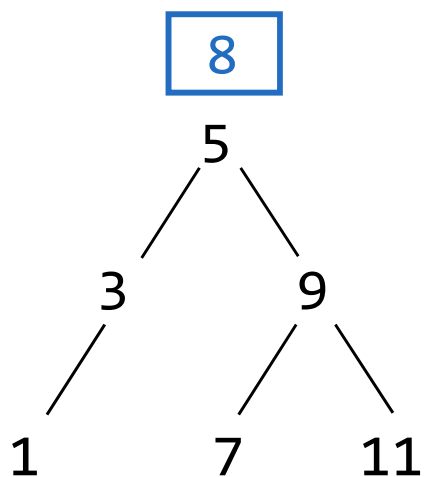
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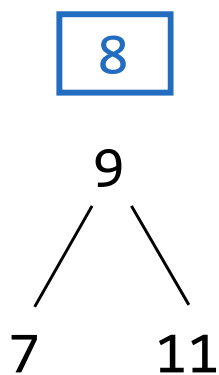
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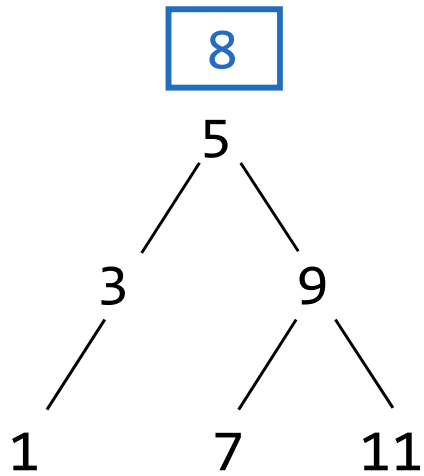


Left!

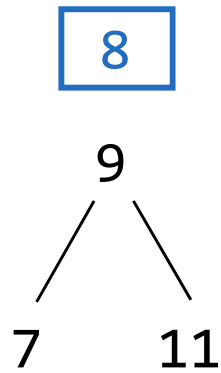




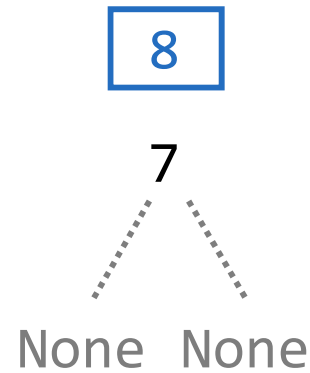
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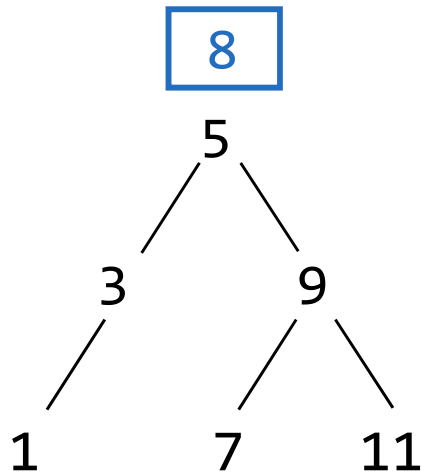
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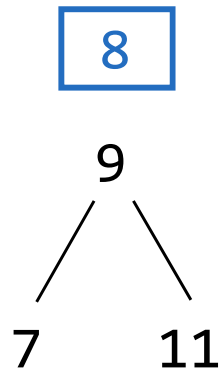
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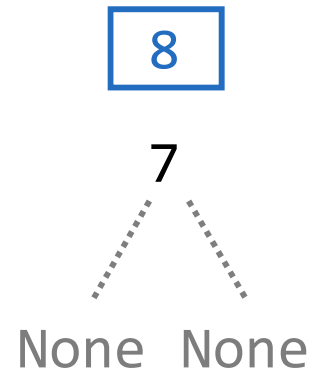
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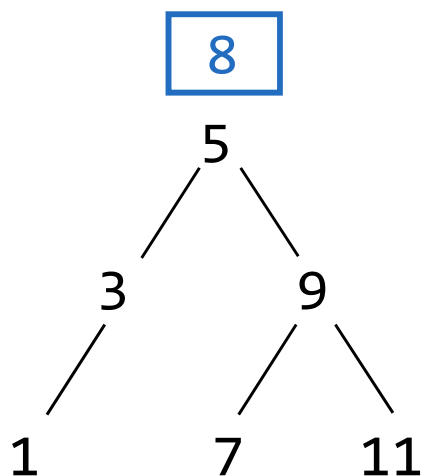
Left!



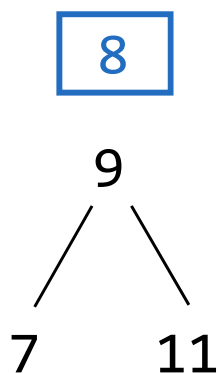
Right!



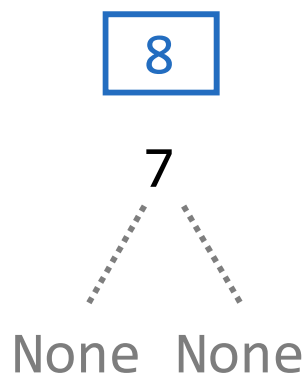
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Right!



Left!



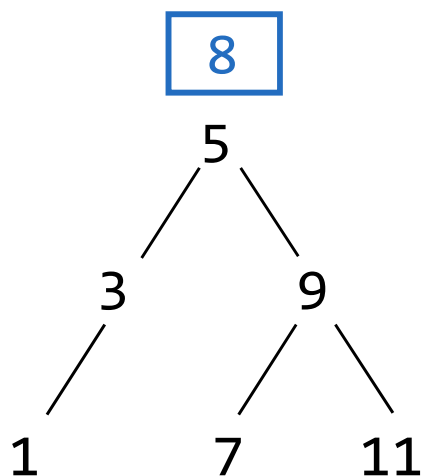
Right!



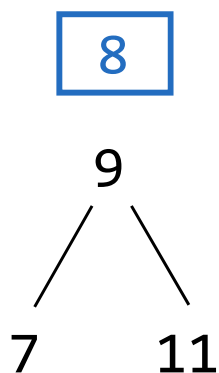
None



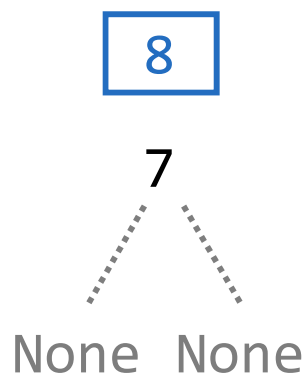
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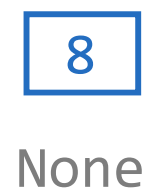
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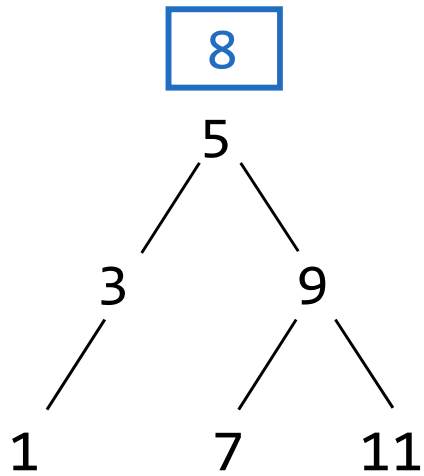
Right!



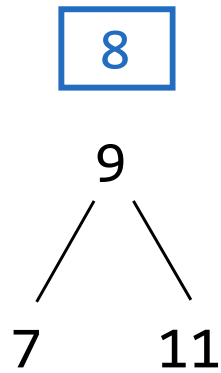
Stop!



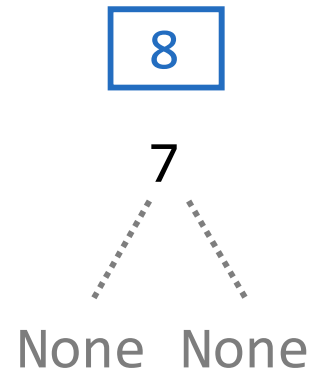
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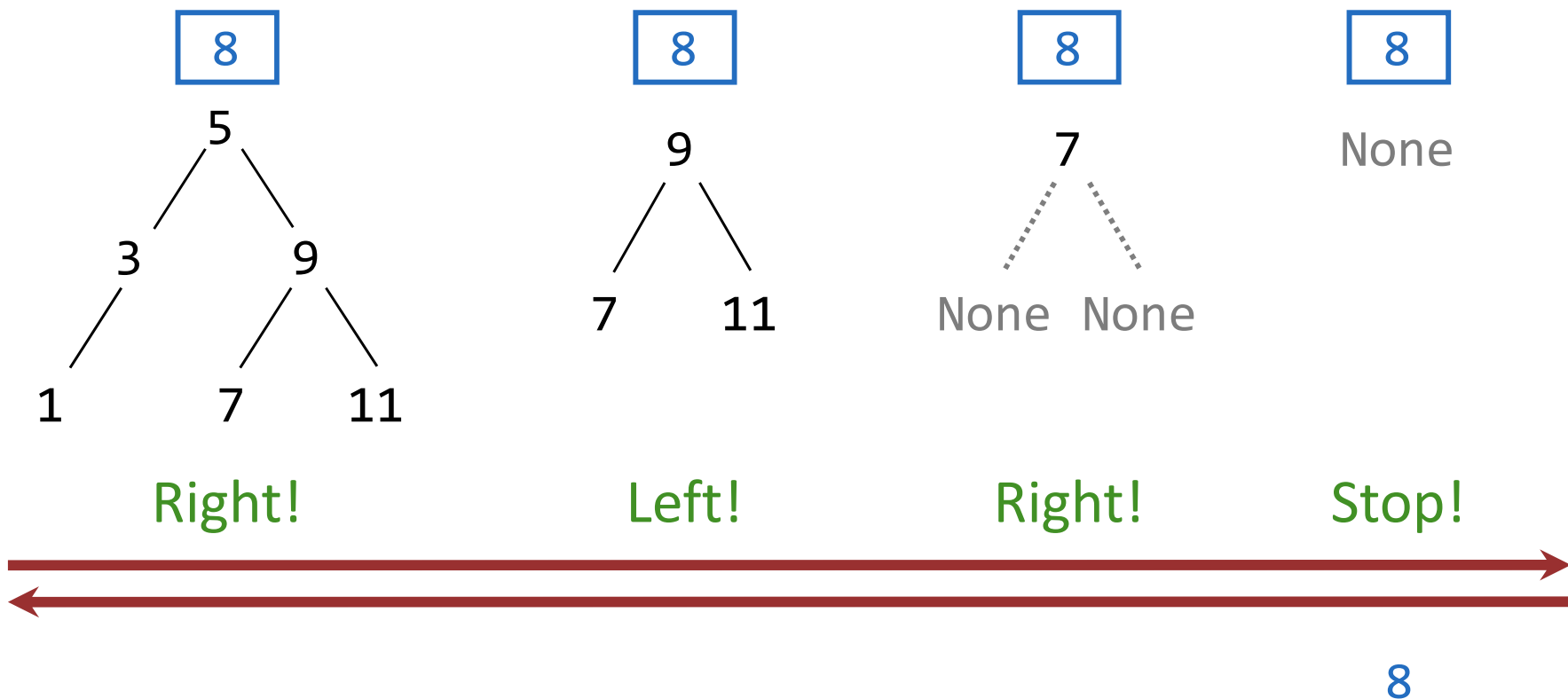
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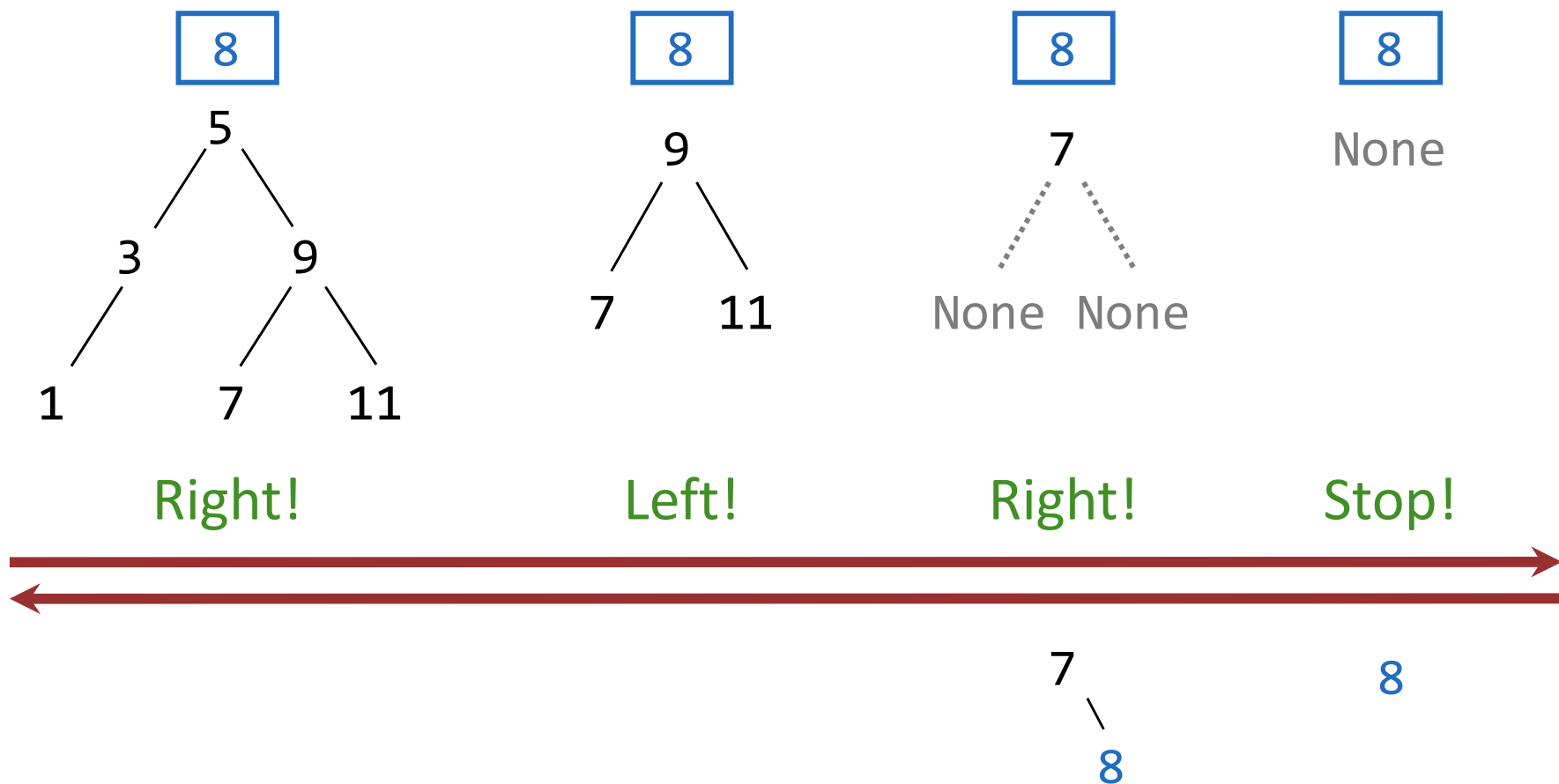
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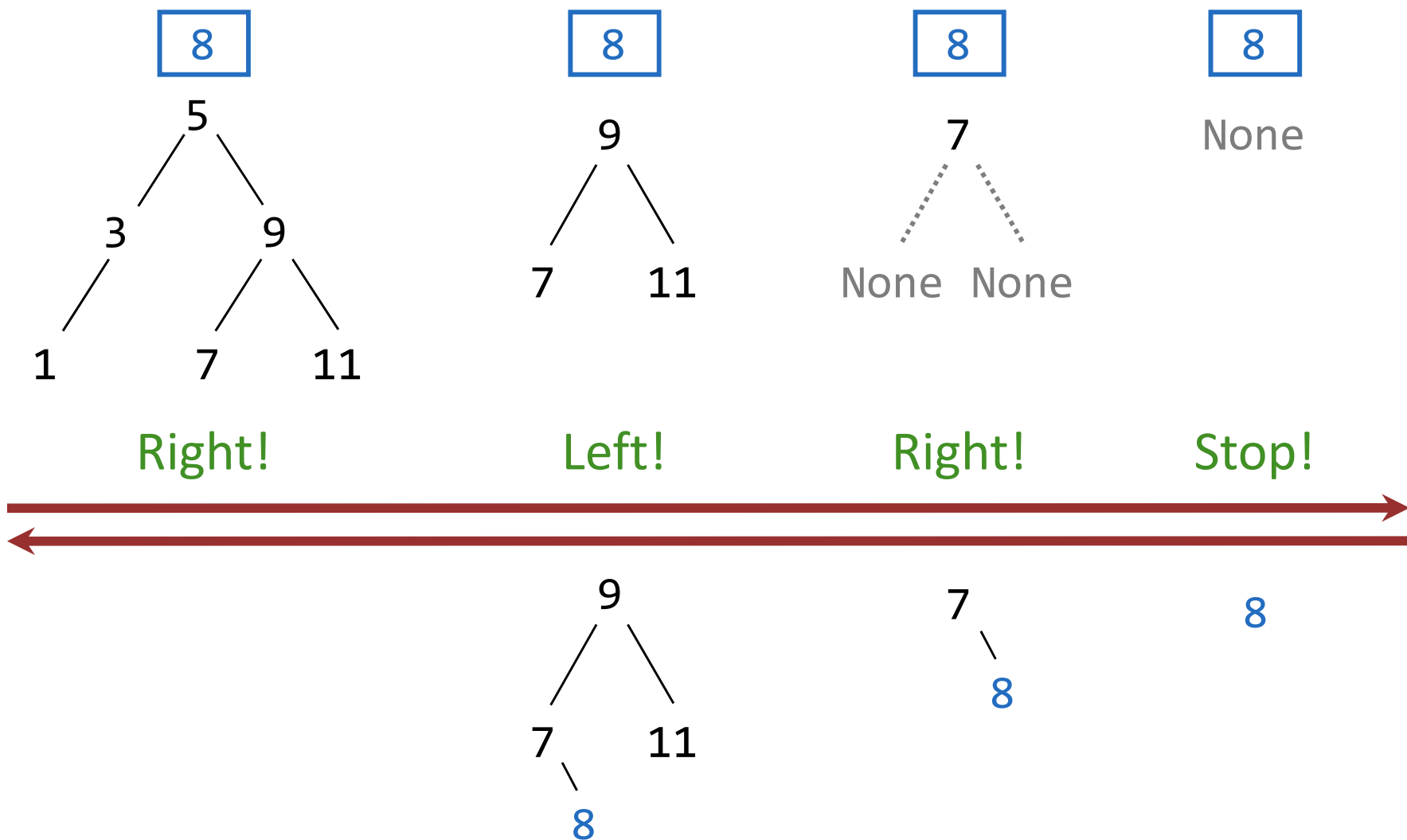
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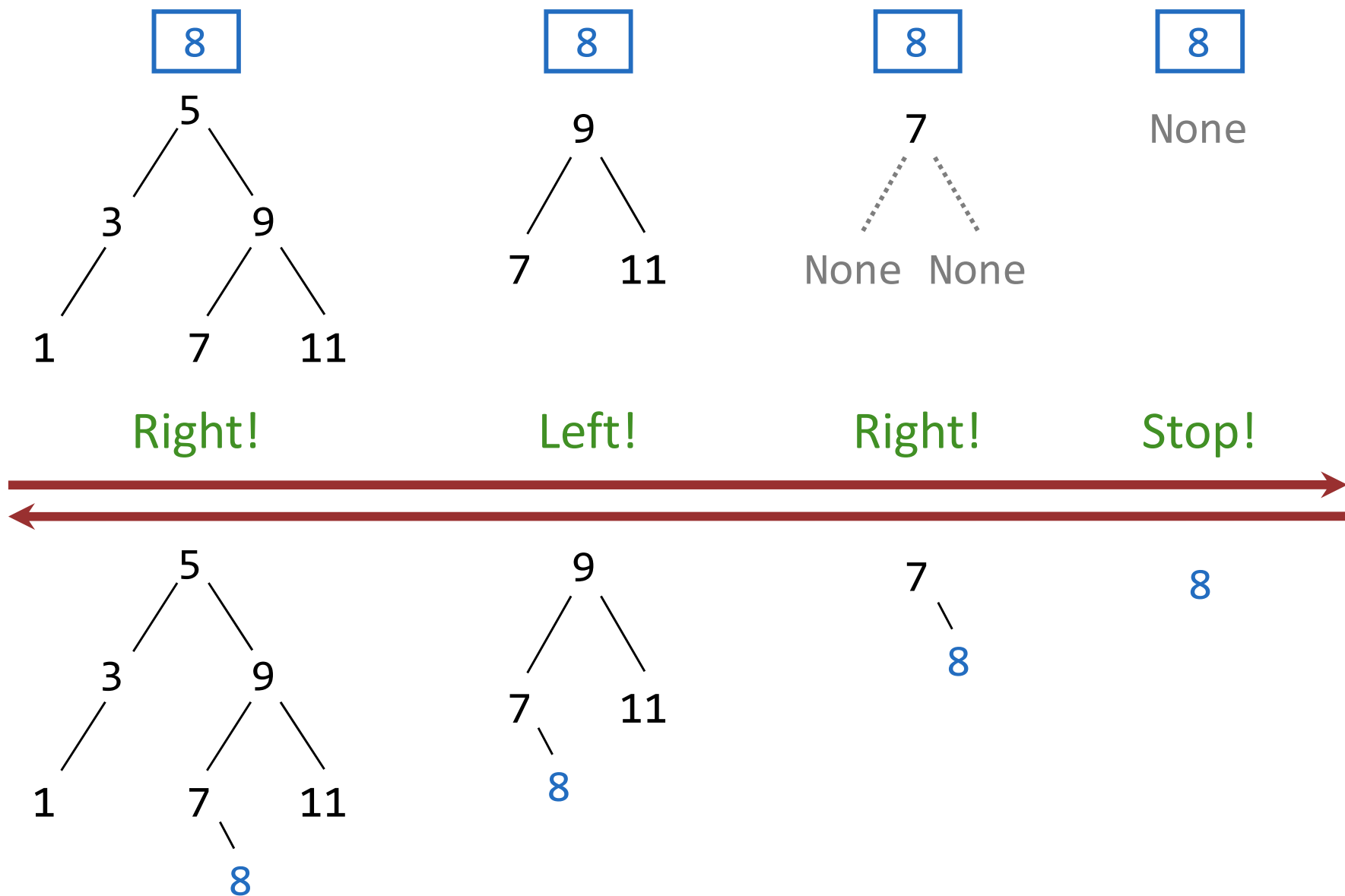


# Adjoining to a Tree Set





# Adjoining to a Tree Set



# What Did I Leave Out?



# What Did I Leave Out?



Sets as ordered sequences:

# What Did I Leave Out?



Sets as ordered sequences:

- Adjoining an element to a set

# What Did I Leave Out?



Sets as ordered sequences:

- Adjoining an element to a set
- Union of two sets

# What Did I Leave Out?



Sets as ordered sequences:

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Sets as binary trees:

# What Did I Leave Out?



Sets as ordered sequences:

- Adjoining an element to a set
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Sets as binary trees:

- Intersection of two sets

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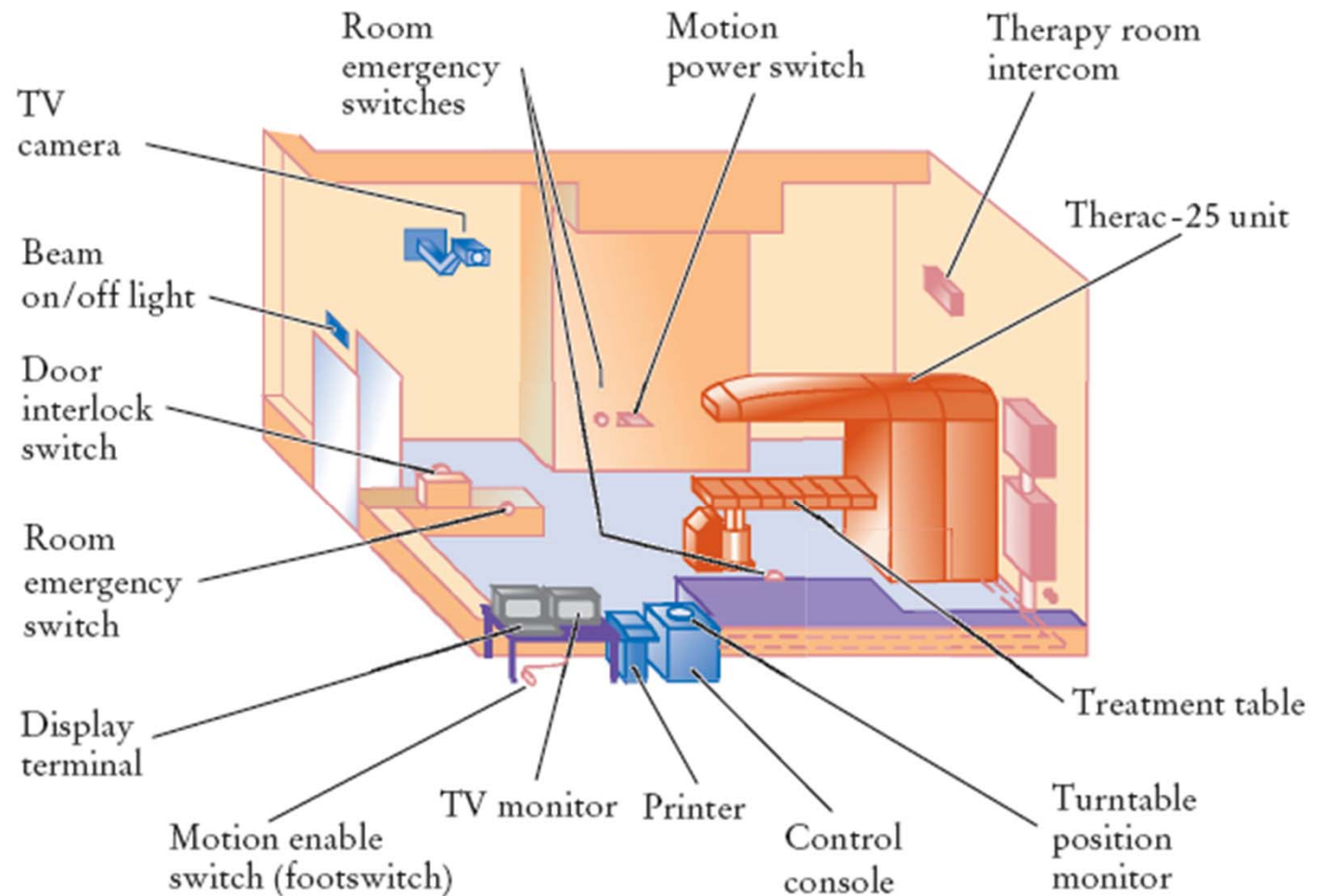
That's homework 9!

- Why things go wrong
- What can we do about this

# Therac-25 Case Study



# Therac-25 Case Study

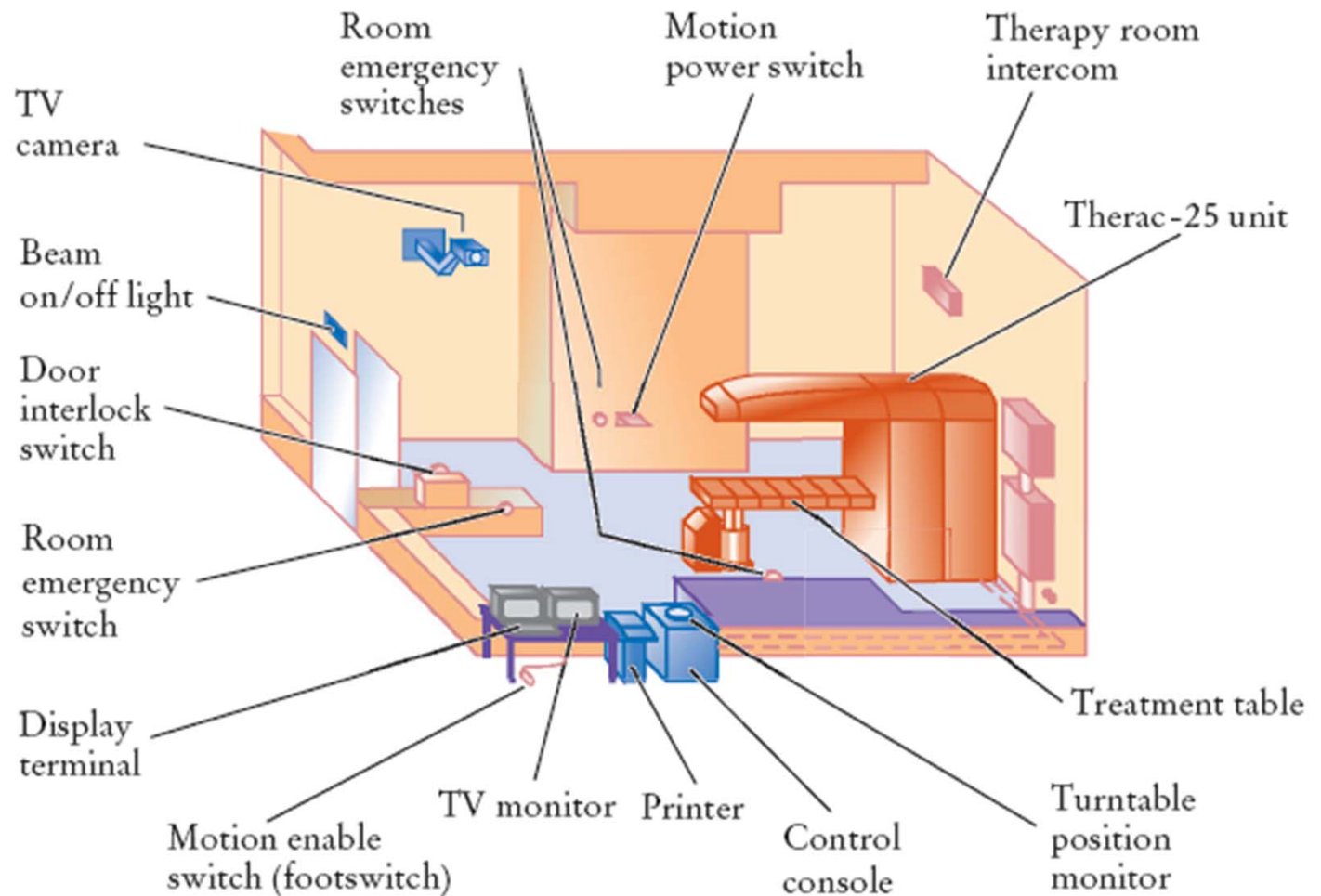


**Figure 9** Typical Therac-25 Facility

# Therac-25 Case Study



## □ Medical imaging device



**Figure 9** Typical Therac-25 Facility

# Therac-25 Case Study



# Therac-25 Case Study



- What happened?

# Therac-25 Case Study



- What happened?
- 6 serious injuries



# Therac-25 Case Study



- What happened?
- 6 serious injuries
- 4 deaths

# Therac-25 Case Study



- What happened?
- 6 serious injuries
- 4 deaths
- Otherwise effective – saved hundreds of lives

# Lesson to be learned



# Lesson to be learned



- Social responsibility in engineering

# Lesson to be learned



- Social responsibility in engineering
- First real incident of fatal software failure

# Lesson to be learned



- Social responsibility in engineering
- First real incident of fatal software failure
- Bigger issue
  - No bad guys
  - Honestly believed there was nothing wrong

# “Software Rot”



- Other engineering fields: clear sense of degradation and decay
- Can software become brittle or fractured?

# A bigger picture





# A bigger picture



- **All software is part of a bigger system**

- **All software is part of a bigger system**
- Software degrades because:
  - Other piece of software changes
  - Hardware changes
  - Environment changes

# Ex: Compatibility Issues



# A bigger issue

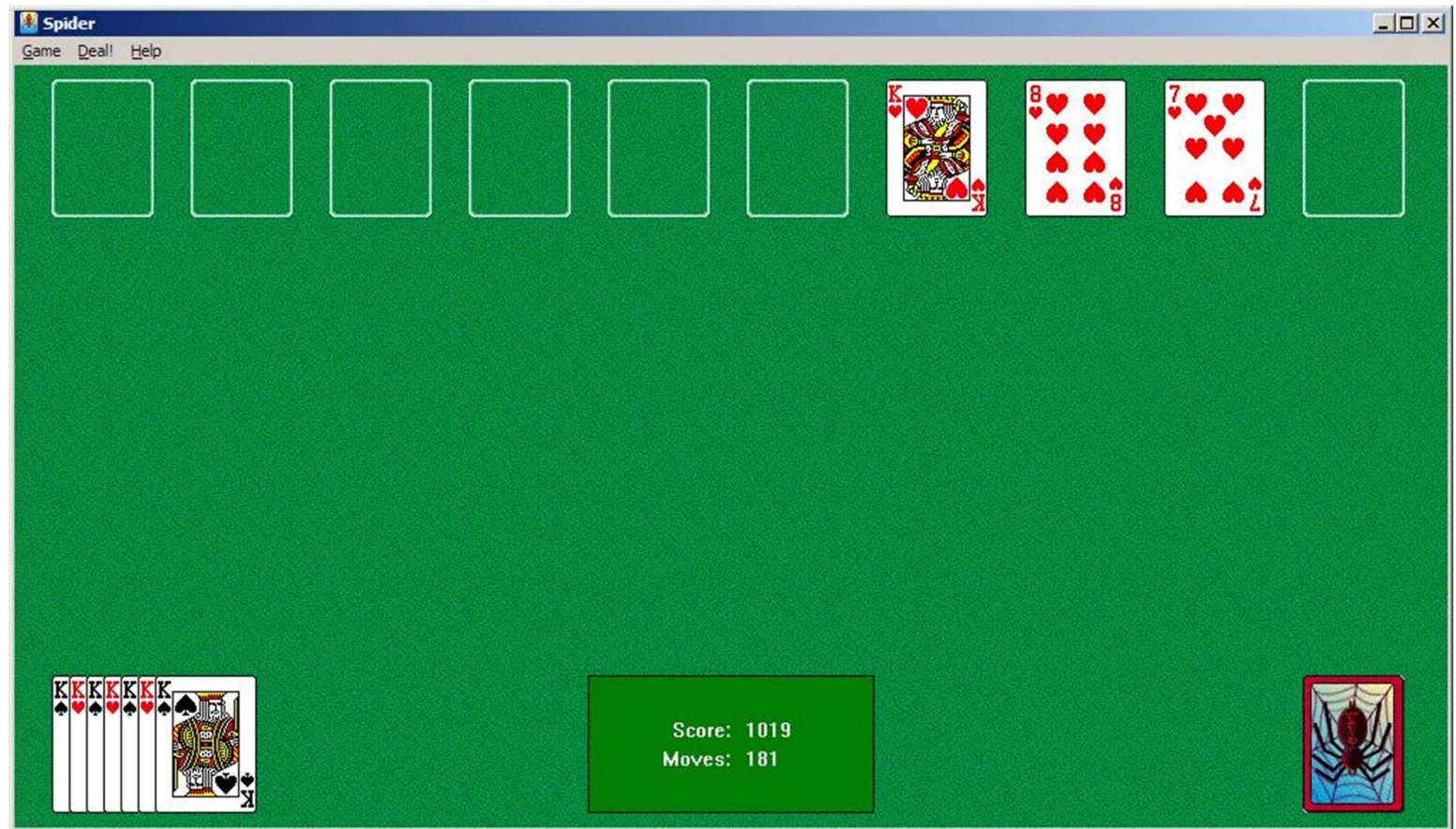


- The makers of the Therac did not fully understand the **complexity** of their software
- Complexity of constructs in other fields more apparent

# A “simple” program



# A “simple” program



# A “simple” program



- This program can delete any file you can

# Complexity in the Therac-25





# Complexity in the Therac-25



- Abundant user interface issues

# Complexity in the Therac-25



- Abundant user interface issues
- Cursor position and field entry

# Complexity in the Therac-25



- Abundant user interface issues
- Cursor position and field entry
- Default values

# Complexity in the Therac-25



- Abundant user interface issues
- Cursor position and field entry
- Default values
- Too many error messages

# Too many error messages



# Too many error messages



# (More) Complexity in the Therac-25



- ❑ No atomic test-and-set
- ❑ No hardware interlocks

# How can we solve these things?



- Know your user
- Fail-Soft (or Fail-Safe)
- Audit Trail
- Correctness from the start
- Redundancy



# Fail-Soft (or Fail-Safe)



```
def mutable_rlist():
    def dispatch(message, value=None):
        nonlocal contents
        if message == 'first':
            return first(contents)
        if message == 'rest':
            return rest(contents)
        if message == 'len':
            return len_rlist(contents)
        ...

    return dispatch
```

# Fail-Soft (or Fail-Safe)



```
def mutable_rlist():
    def dispatch(message, value=None):
        nonlocal contents
        if message == 'first':
            return first(contents)
        if message == 'rest':
            return rest(contents)
        if message == 'len':
            return len_rlist(contents)
        ...
        else:
            print('Unknown message')
    return dispatch
```

# Correctness from the start



- Edsger Dijkstra: “On the Cruelty of Really Teaching Computing Sciences”
- CS students shouldn’t use computers
- Rigorously prove correctness of their programs
  
- Correctness proofs
- Compilation (pre-execution) analysis

# On debugging



- ❑ Black box debugging
- ❑ Glass box debugging
- ❑ Don't break what works
  
- ❑ Golden rule of debugging...

□ “Debug by subtraction, not by addition”

□ Prof. Brian Harvey