



CS61A Lecture 25

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Announcements



- HW8 due tonight at 7pm

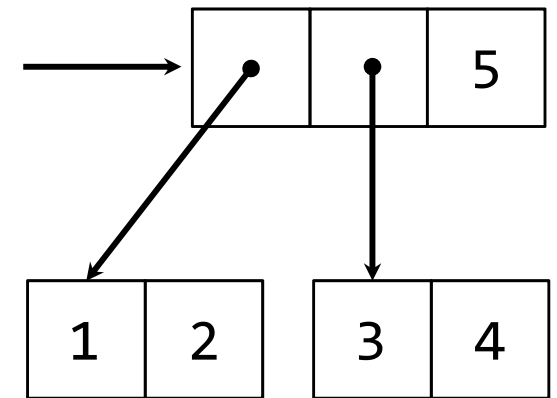
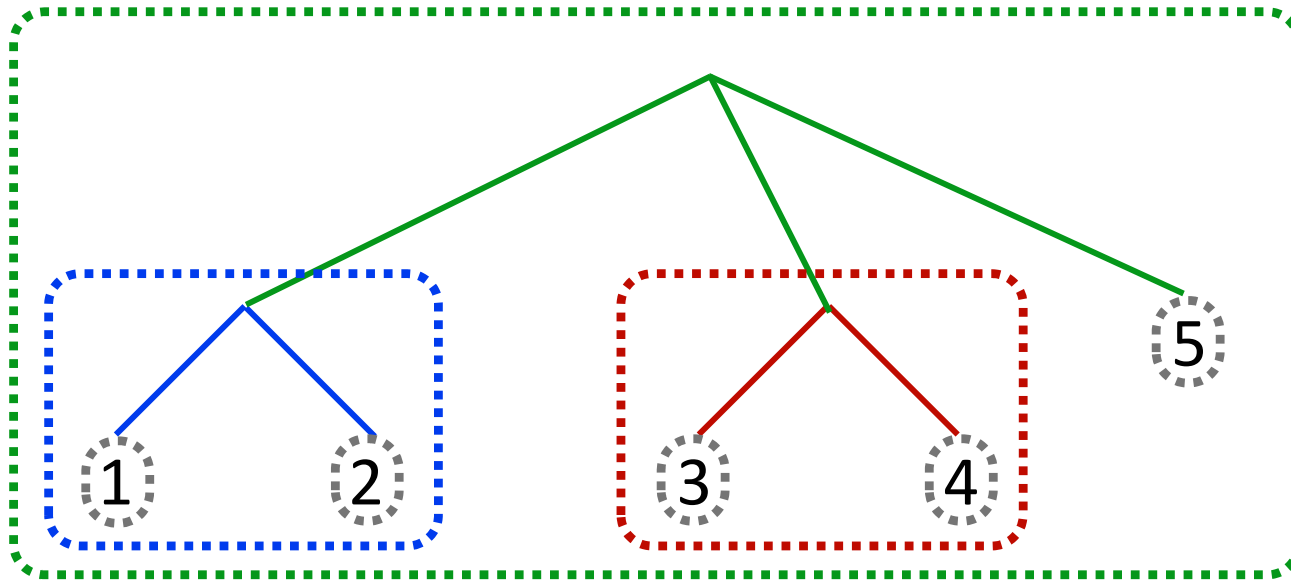
- Midterm 2 Thursday at 7pm
 - See course website for more information

Tree Structured Data



Nested Sequences are Hierarchical Structures.

$((1, 2), (3, 4), 5)$



In every tree, a vast forest

Example: <http://goo.gl/0h6n5>

Recursive Tree Processing



Tree operations typically make recursive calls on branches

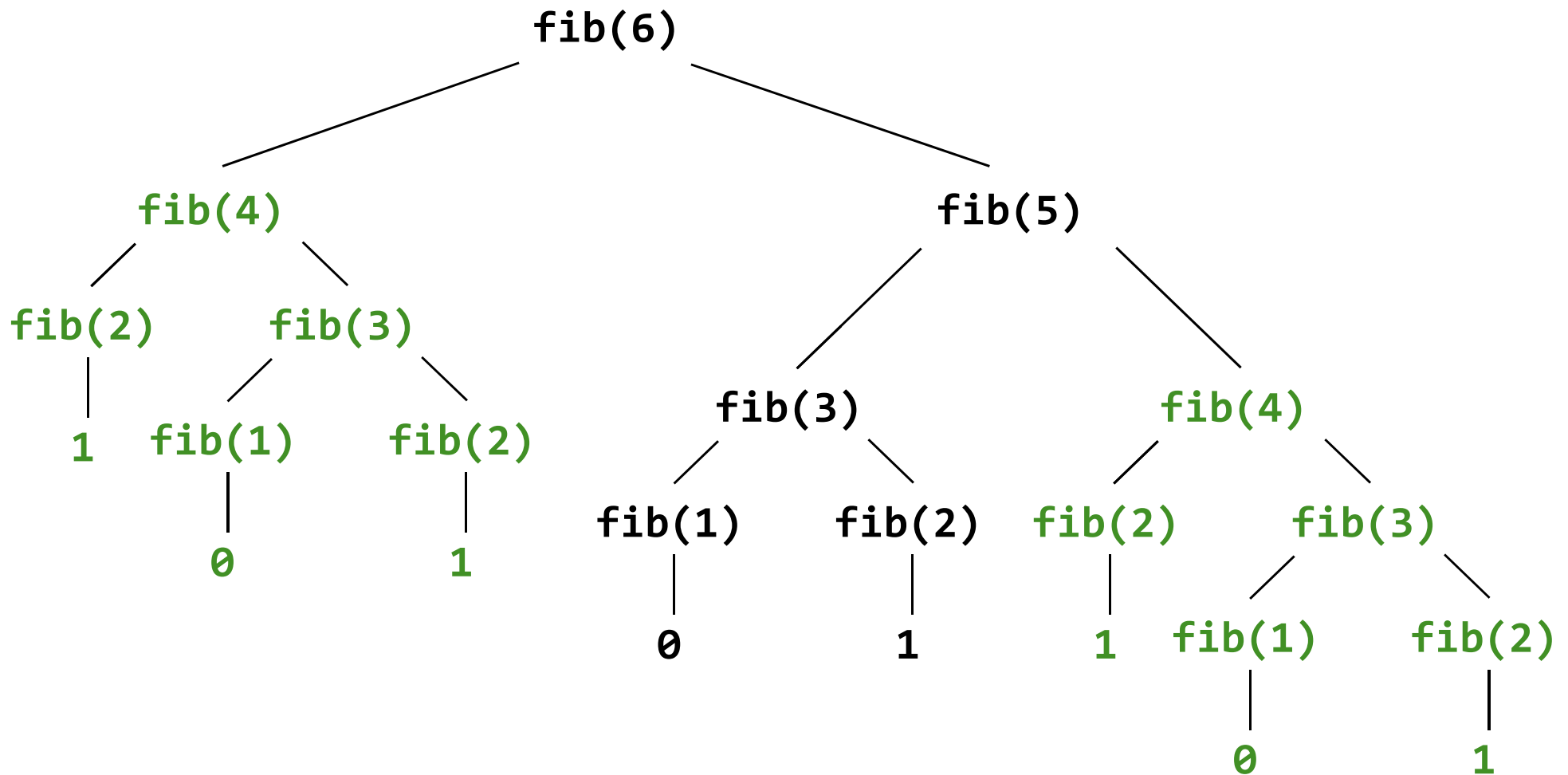
```
def count_leaves(tree):  
    if type(tree) != tuple:  
        return 1  
    return sum(map(count_leaves, tree))
```

```
def map_tree(tree, fn):  
    if type(tree) != tuple:  
        return fn(tree)  
    return tuple(map_tree(branch, fn)  
                  for branch in tree)
```

Trees with Internal Node Values



Trees can have values at internal nodes as well as their leaves.



Trees with Internal Node Values



Trees can have values at internal nodes as well as their leaves.

```
class Tree(object):
    def __init__(self, entry, left=None, right=None):
        self.entry = entry
        self.left = left
        self.right = right

def fib_tree(n):
    if n == 1:
        return Tree(0)
    if n == 2:
        return Tree(1)
    left = fib_tree(n - 2)
    right = fib_tree(n - 1)
    return Tree(left.entry + right.entry, left, right)
```

Memoization



Tree recursive functions can compute the same thing many times

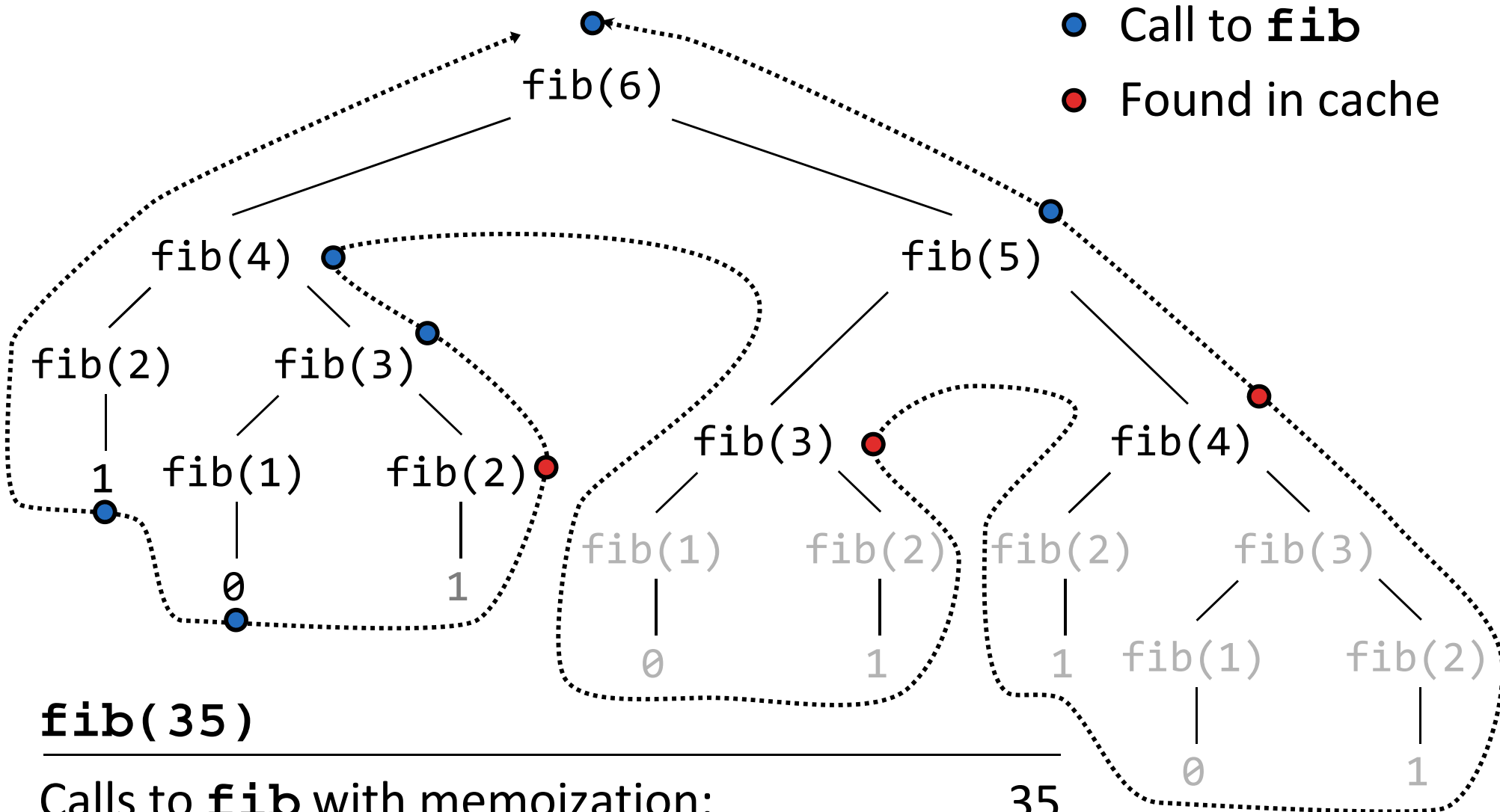
Idea: Remember the results that have been computed before

```
def memo(f):  
    cache = {}  
    def memoized(n):  
        if n not in cache:  
            cache[n] = f(n)  
        return cache[n]  
    return memoized
```

Keys are arguments that map to return values

Same behavior as f , if f is a pure function

Memoized Tree Recursion



fib(35)

Calls to **fib** with memoization: 35

Calls to **fib** without memoization: 18,454,929

Orders of Growth



Iterative, recursive, and memoized implementations are not the same.

	<u>Time</u>	<u>Space</u>
<pre>def fib_iter(n): prev, curr = 1, 0 for _ in range(n - 1): prev, curr = curr, prev + curr return curr</pre>	$\Theta(n)$	$\Theta(1)$
<pre>def fib(n): if n == 1: return 0 if n == 2: return 1 return fib(n - 2) + fib(n - 1)</pre>	$\Theta(\phi^n)$	$\Theta(n)$
<pre>fib = memo(fib)</pre>	$\Theta(n)$	$\Theta(n)$

Sets



One more built-in Python container type

- Set literals are enclosed in braces
- Duplicate elements are removed on construction
- Sets are unordered, just like dictionary entries

```
>>> s = {3, 2, 1, 4, 4}
```

```
>>> s
```

```
{1, 2, 3, 4}
```

```
>>> 3 in s
```

```
True
```

```
>>> len(s)
```

```
4
```

```
>>> s.union({1, 5})
```

```
{1, 2, 3, 4, 5}
```

```
>>> s.intersection({6, 5, 4, 3})
```

```
{3, 4}
```

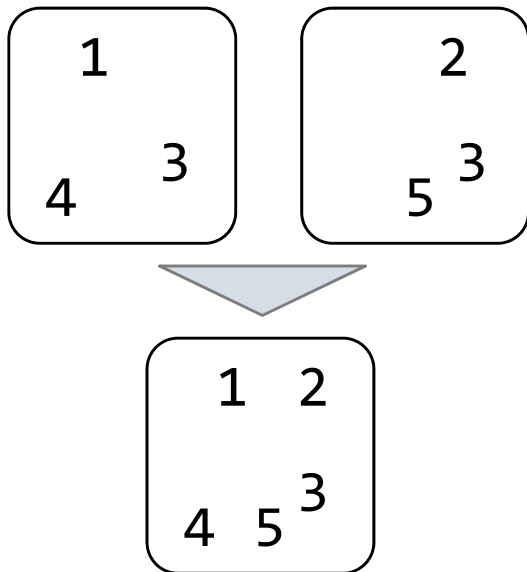
Implementing Sets



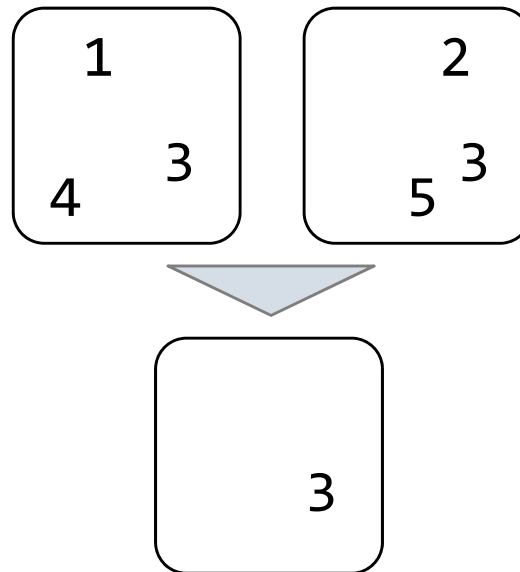
What we should be able to do with a set:

- Membership testing: Is a value an element of a set?
- Union: Return a set with all elements in *set1* **or** *set2*
- Intersection: Return a set with any elements in *set1* **and** *set2*
- Adjunction: Return a set with all elements in *s* and a value *v*

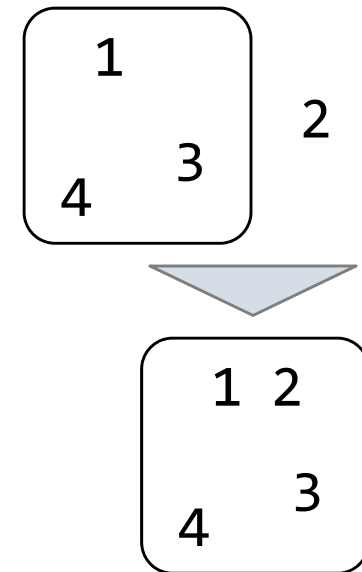
Union



Intersection



Adjunction



Sets as Unordered Sequences



Proposal 1: A set is represented by a recursive list that contains no duplicate items

This is how we implemented dictionaries

```
def empty(s):  
    return s is Rlist.empty  
  
def set_contains(s, v):  
    if empty(s):  
        return False  
    elif s.first == v:  
        return True  
    return set_contains(s.rest, v)
```

Sets as Unordered Sequences



Time order of growth

```
def adjoin_set(s, v):  
    if set_contains(s, v):  
        return s  
    return Rlist(v, s)
```

$$\Theta(n)$$

The size of
the set

```
def intersect_set(set1, set2):  
    f = lambda v: set_contains(set2, v)  
    return filter_rlist(set1, f)
```

$$\Theta(n^2)$$

Assume sets are
the same size

```
def union_set(set1, set2):  
    f = lambda v: not set_contains(set2, v)  
    set1_not_set2 = filter_rlist(set1, f)  
    return extend_rlist(set1_not_set2, set2)
```

$$\Theta(n^2)$$